Unit Disk Representations

of Embedded Trees, Outerplanar and Multi-Legged Graphs

Sujoy Bhore, Maarten Löffler, Soeren Nickel and Martin Nöllenburg 16.09.2021 · GD 2021

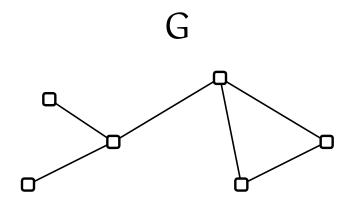






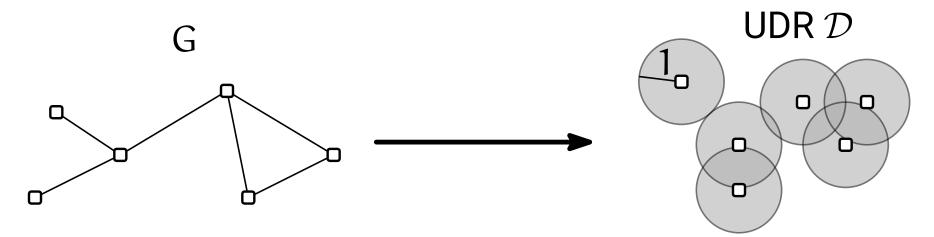


A graph G is a unit disk graph



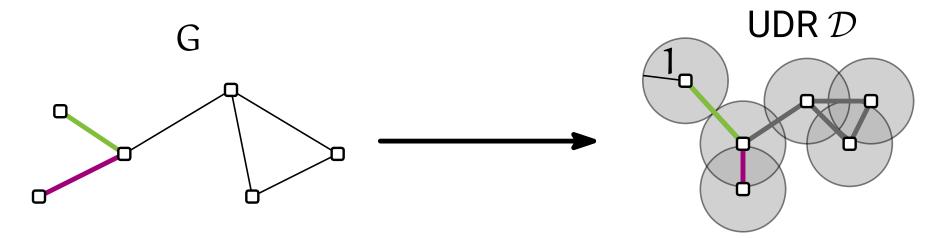


A graph G is a unit disk graph if it has a unit disk representation \mathcal{D} .





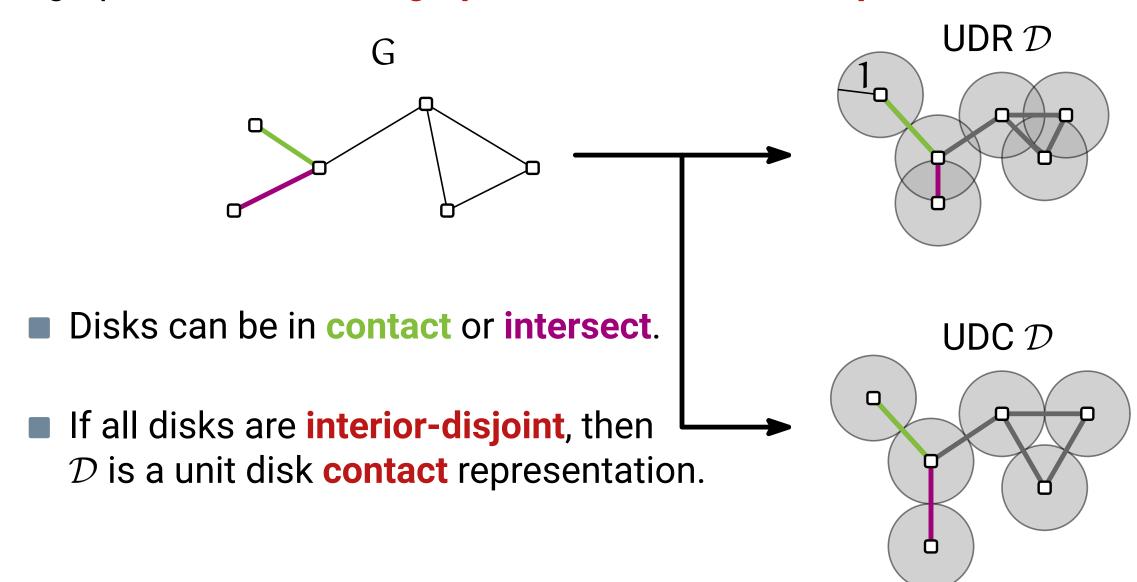
A graph G is a unit disk graph if it has a unit disk representation \mathcal{D} .



Disks can be in contact or intersect.

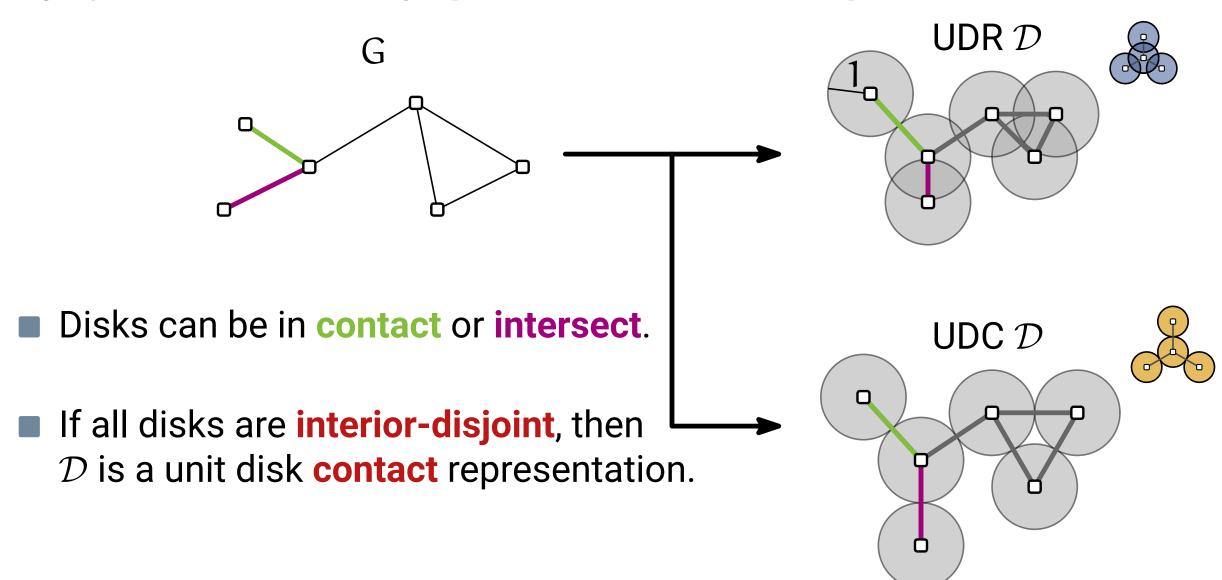


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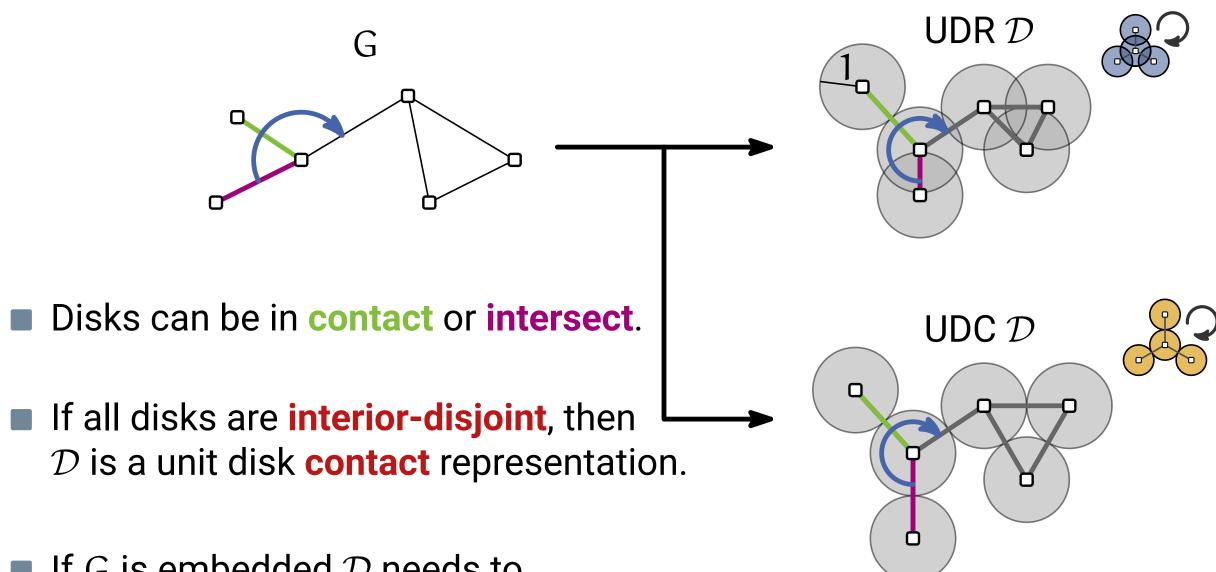


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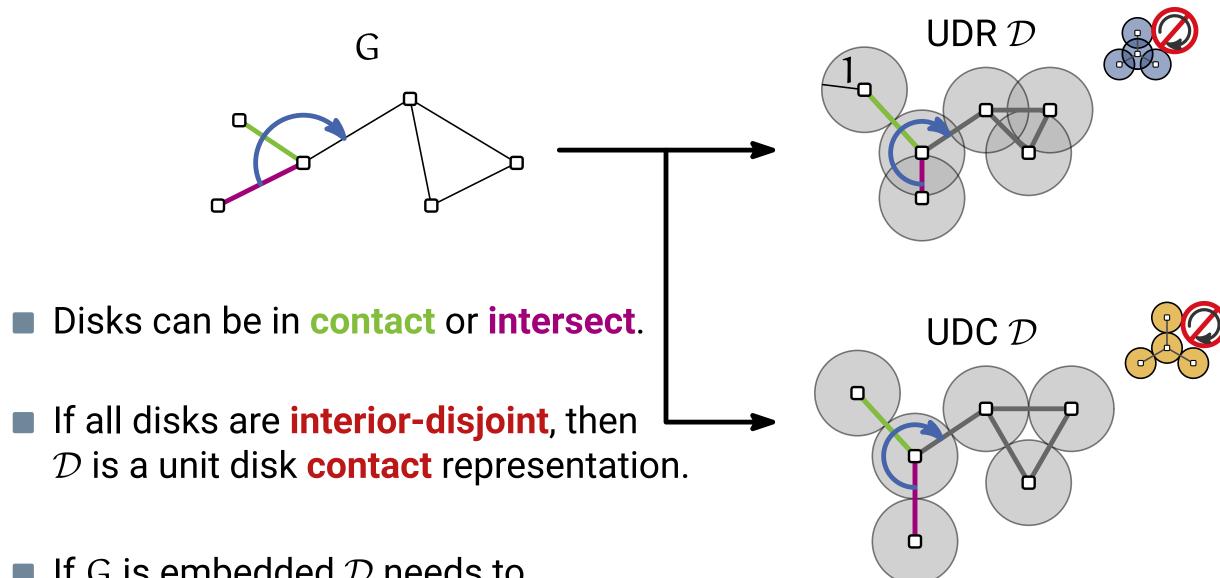
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■ If G is embedded \mathcal{D} needs to induce the same rotation system.



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Recognition of unit disk (contact) graphs is NP-hard for planar graphs.¹

¹[Breu & Kirkpatrick, 1998]



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		planar	outerplanar	tree	caterpillar
contact gra	phs	NP-hard ¹			
embed					
disk gra	phs	NP-hard ¹			
embed	lded				

¹[Breu & Kirkpatrick, 1998]

3 - 2



Recognition of unit disk (contact) graphs is NP-hard for planar graphs.¹

•	, , , ,			
	planar	outerplanar	tree	caterpillar
contact graphs	NP-hard ¹	NP-hard ²		$O(n)^2$
embedded				
disk graphs	NP-hard ¹			
embedded				

¹[Breu & Kirkpatrick, 1998] ²[Klemz et al., 2015]



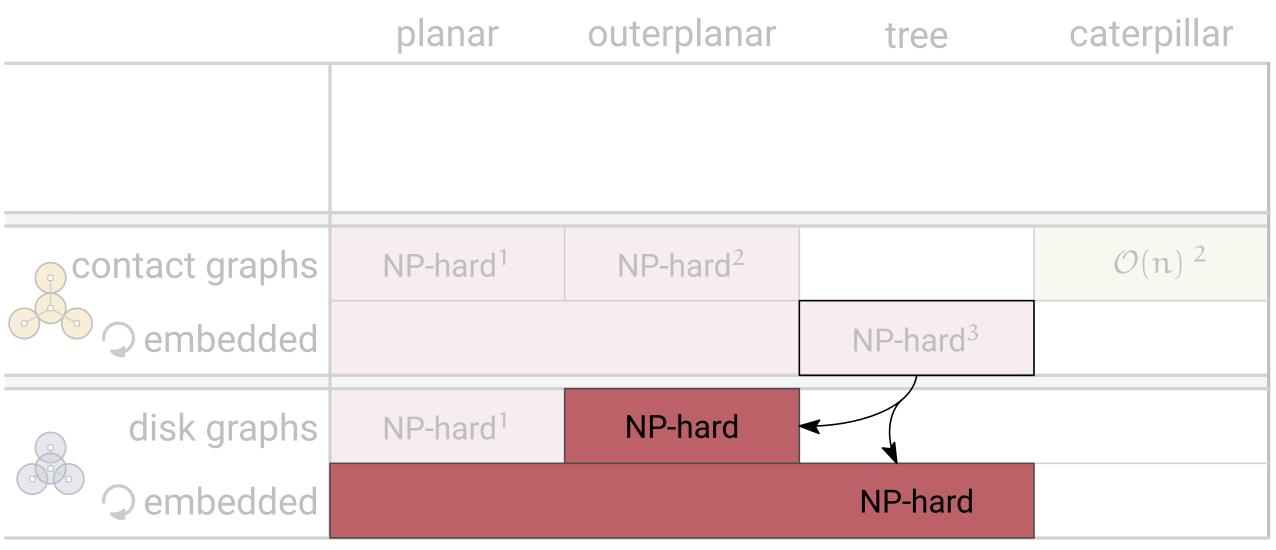
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contact graphs	NP-hard ¹	NP-hard ²		$\mathcal{O}(n)^2$
embedded			NP-hard ³	
disk graphs 2 embedded	NP-hard ¹	NP-hard		$\mathcal{O}(\mathfrak{n})$
			NP-hard	

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3	\		•	5 1
	planar	outerplanar	tree	caterpillar
contact graphs	NP-hard ¹	NP-hard ²	?	$\mathcal{O}(n)^2$
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Recognition of unit disk (contact) graphs is NP-hard for planar graphs. 1

	planar	outerplanar	tree	caterpillar
weak contact			NP-hard ⁴	$\mathcal{O}(\mathfrak{n})^4$
embedded				NP-hard ⁵
© contact graphs	NP-hard ¹	NP-hard ²	?	$\mathcal{O}(n)^2$
embedded			NP-hard ³	?
disk graphs a embedded	NP-hard ¹	NP-hard	?	$\mathcal{O}(n)$
			NP-hard	?

¹[Breu & Kirkpatrick, 1998] ²[Klemz et al., 2015] ⁴[Cleve, 2020]



Soeren Nickel Recognition of Unit Disk Graphs

Recognition of unit disk (contact) graphs is NP-hard for planar graphs.¹

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	planar	outerplanar	tree	caterpillar
weak contact			NP-hard ⁴	(n) § O(n) 4
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embedded			NP-hard	?
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³[Bowen et al., 2015] ⁵[Chiu et al., 2019]

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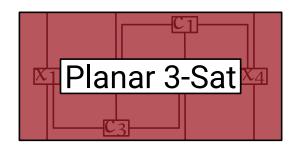
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Reduction from planar 3-Sat

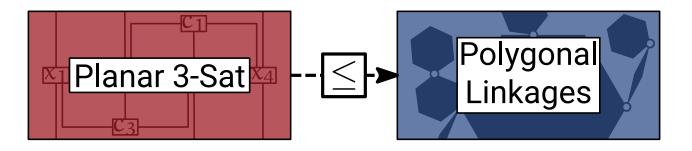
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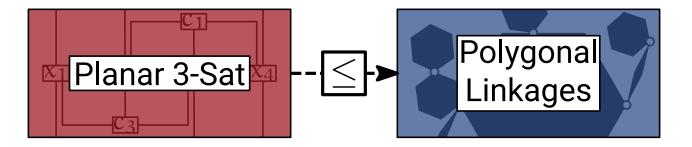
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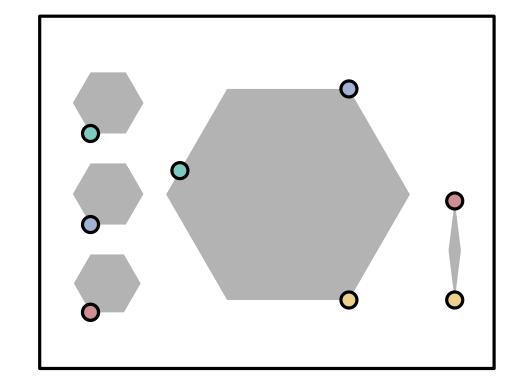




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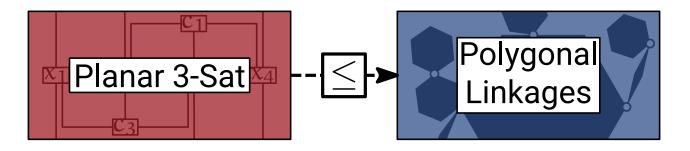


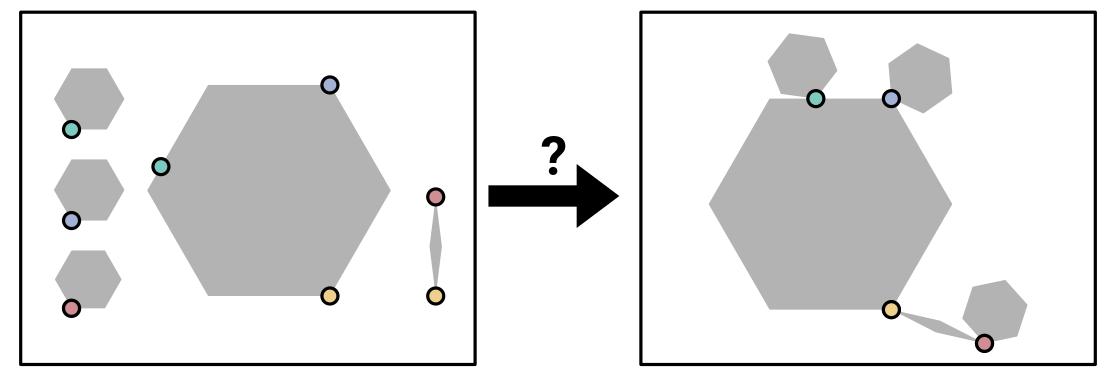




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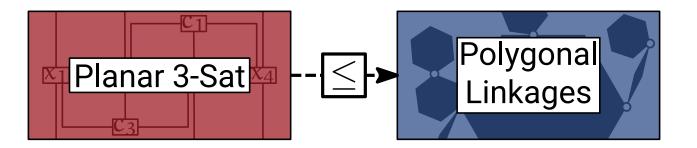


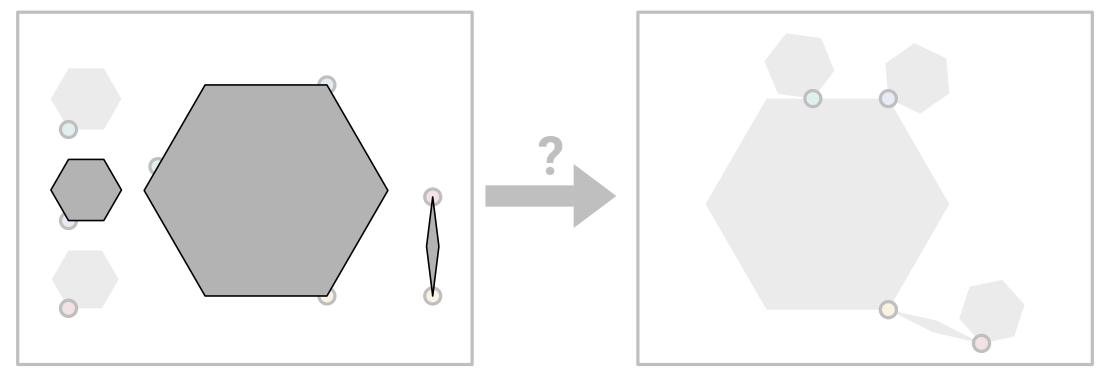


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[Bowen et al., 2015]



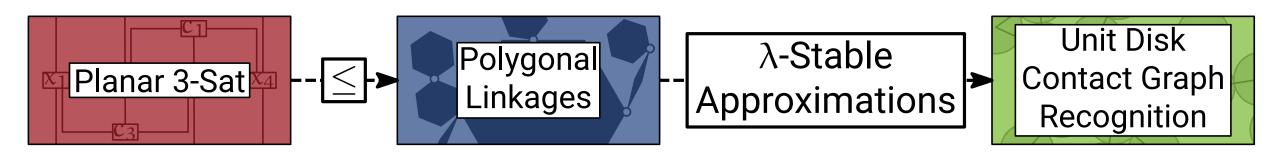


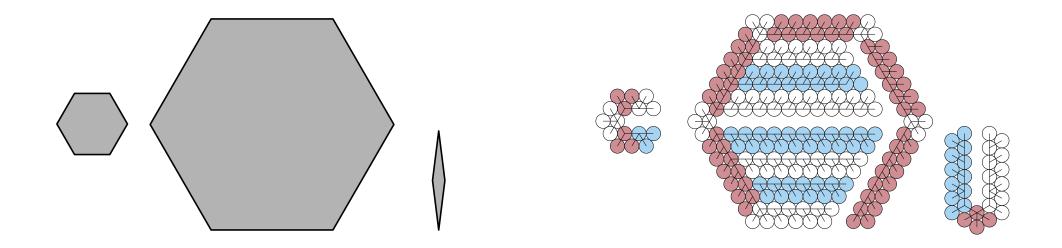


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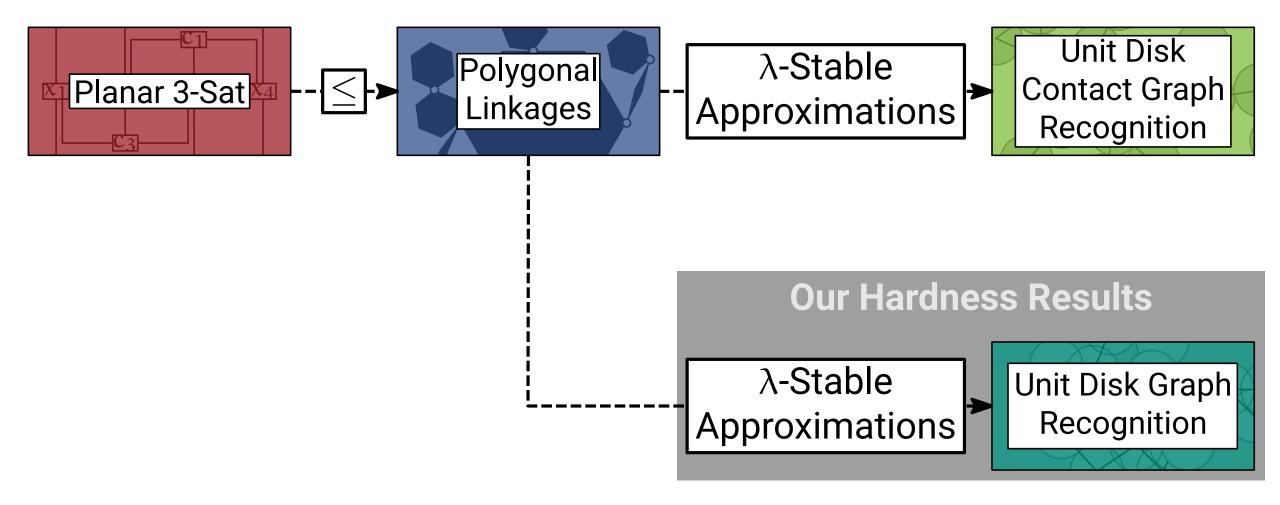






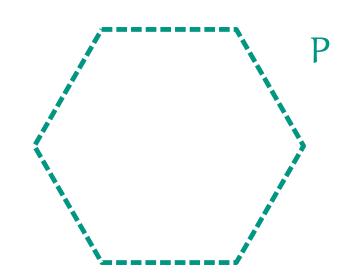
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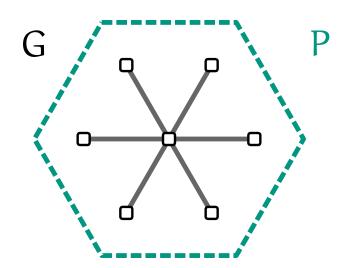


a λ -stable approximation of a polygon P



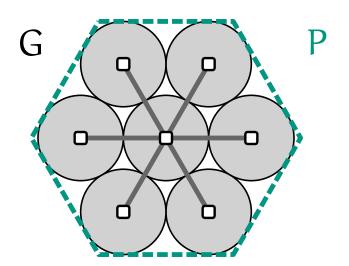


A graph G is a λ -stable approximation of a polygon P



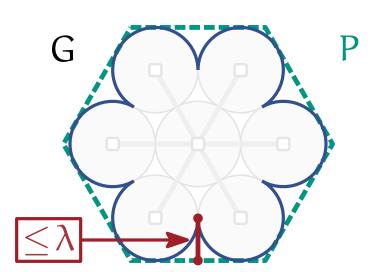


A graph G is a λ -stable approximation of a polygon P, if the union of all disks in every UDC of G





A graph G is a λ -stable approximation of a polygon P, if the union of all disks in every UDC of G has a Hausdorff distance $\leq \lambda$ to P.

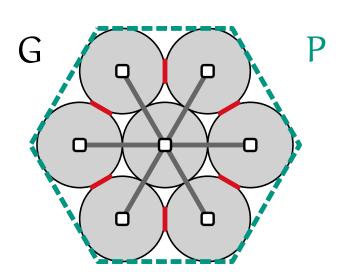


[Bowen et al., 2015]

(directed) Hausdorff-Distance \approx How far do you have to travel from any point on P to reach the UDC boundary of G



A graph G is a λ -stable approximation of a polygon P, if the union of all disks in every UDC of G has a Hausdorff distance $\leq \lambda$ to P.



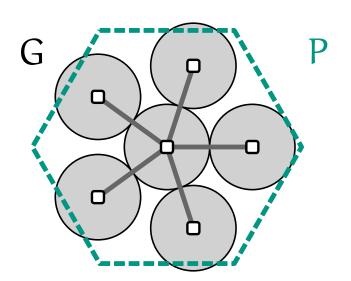
[Bowen et al., 2015]

But:

False adjacencies are not allowed



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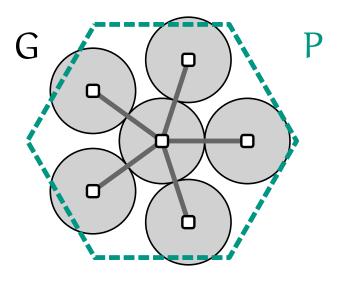
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A graph G is a λ -stable approximation of a polygon P, if the union of all disks in every **UDR** of G has a Hausdorff distance $\leq \lambda$ to P.



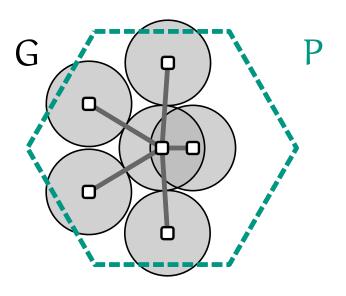
But:

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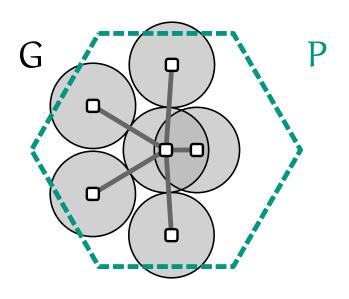
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But:

- False adjacencies are not allowed
- UDR allow for overlap
- Must hold for all UDRs

We have to force a very restricted placement in the UDR.

Bend Restrictions in Embedded UDRs



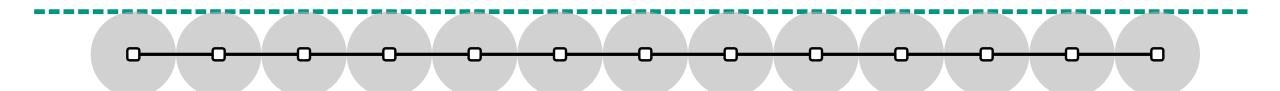
What about tracing a straight line?

Soeren Nickel · Recognition of Unit Disk Graphs

Bend Restrictions in Embedded UDRs

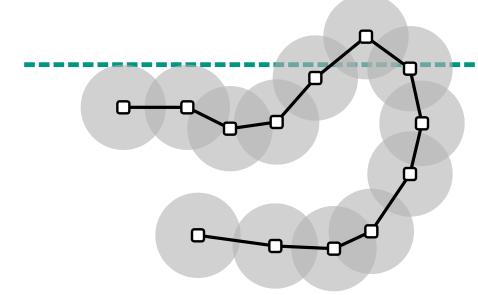


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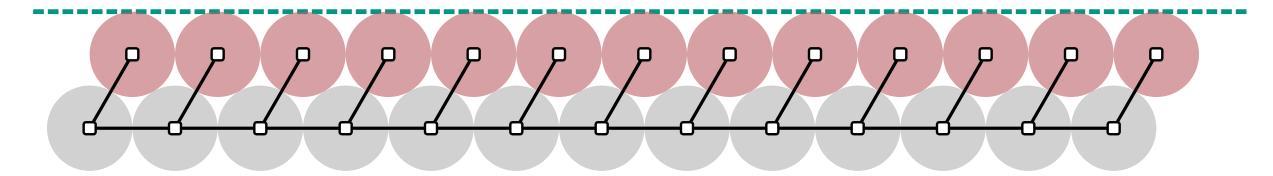


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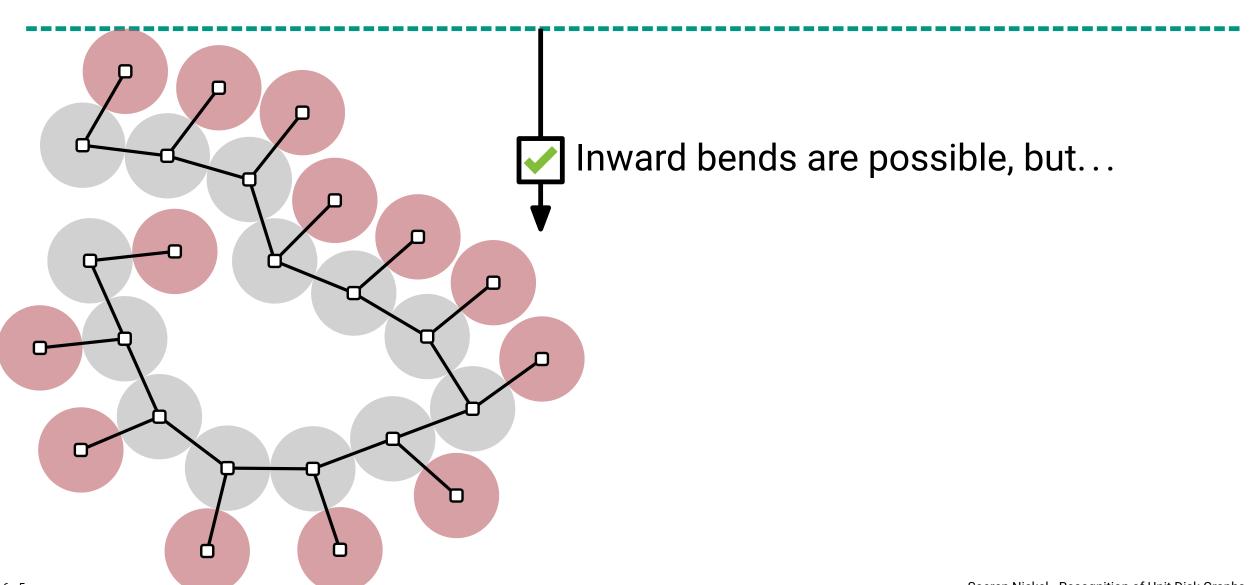


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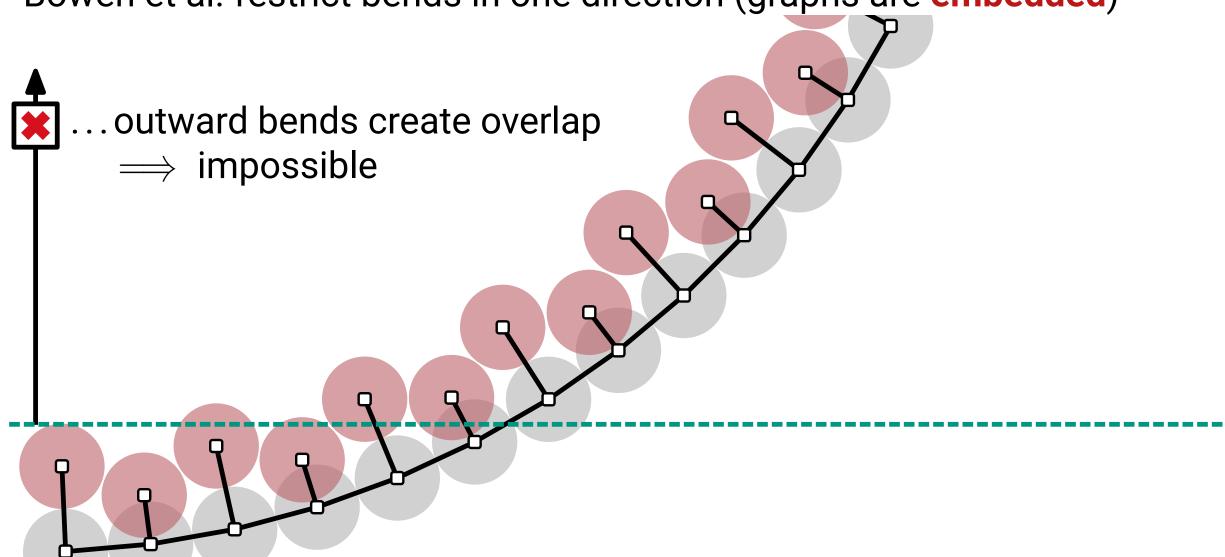


What about tracing a straight line? Bowen et al. restrict bends in one direction (graphs are embedded) . outward bends create overlap impossible



What about tracing a straight line?

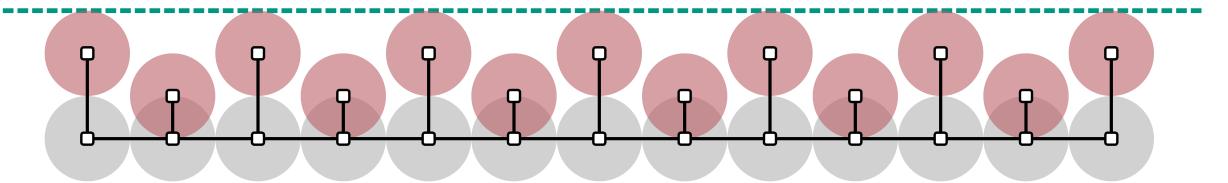
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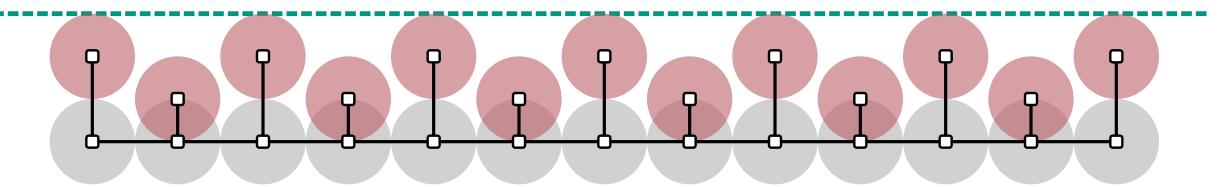


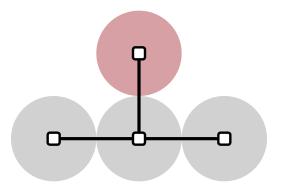
Does not work immediately in a **UDR**Overlaps create space for outward bends



What about tracing a straight line?

Bowen et al. restrict bends in one direction (graphs are embedded)

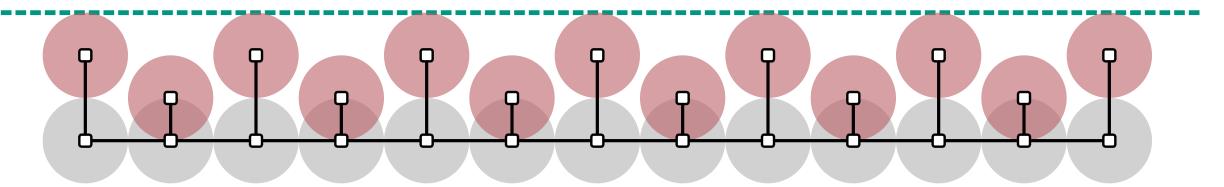


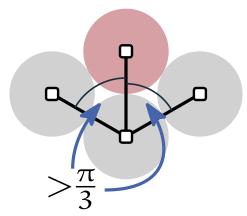




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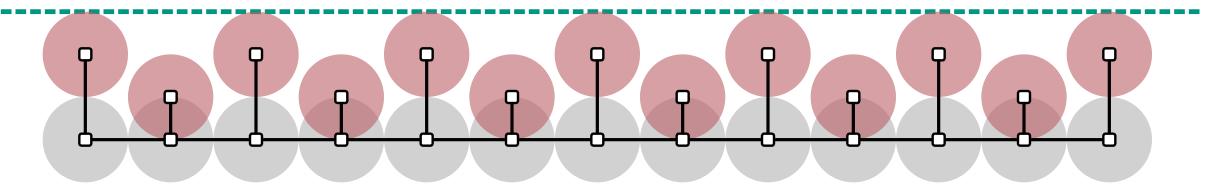


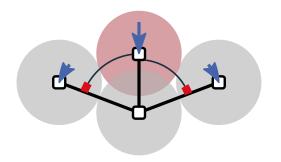




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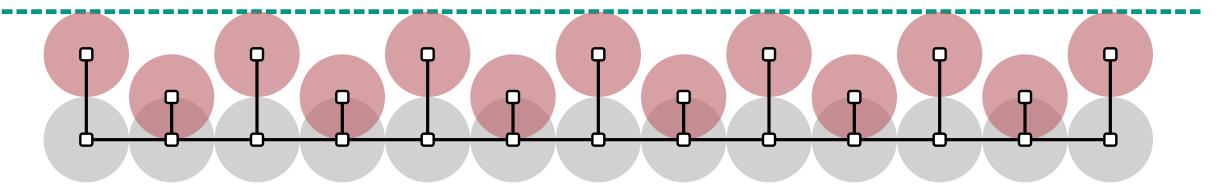


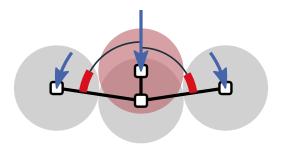




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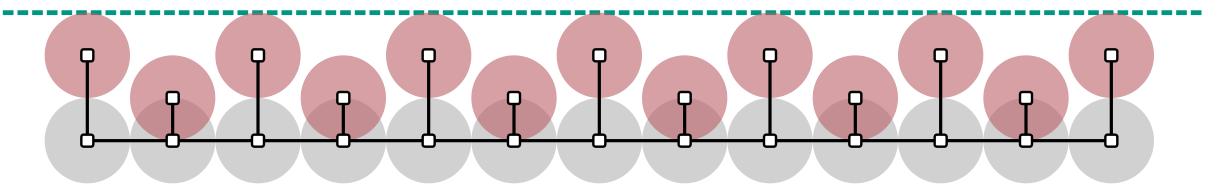


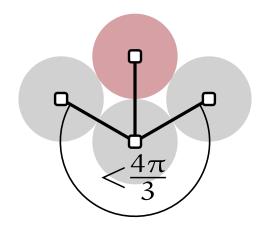




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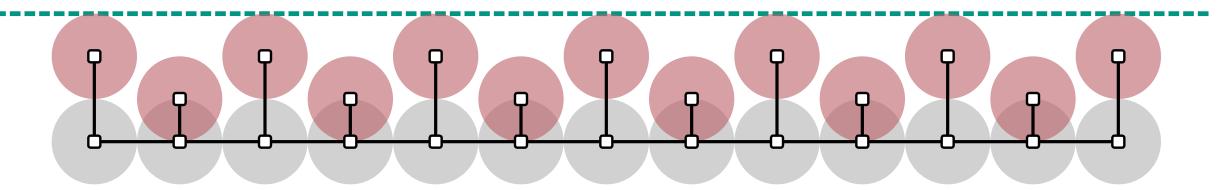


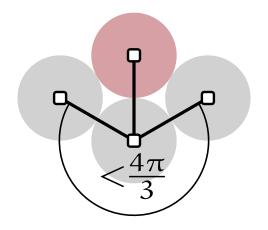


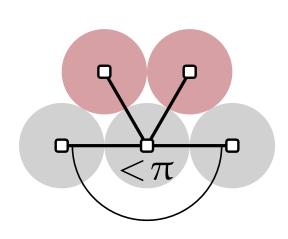


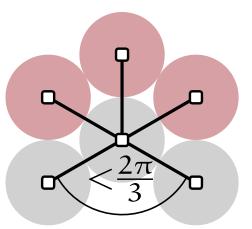
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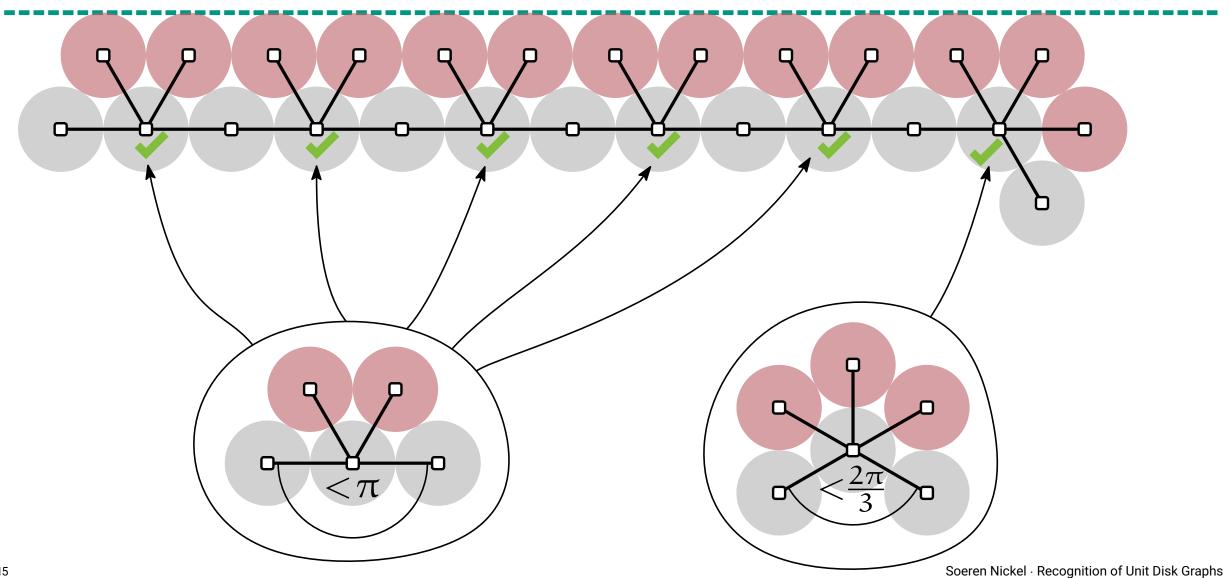






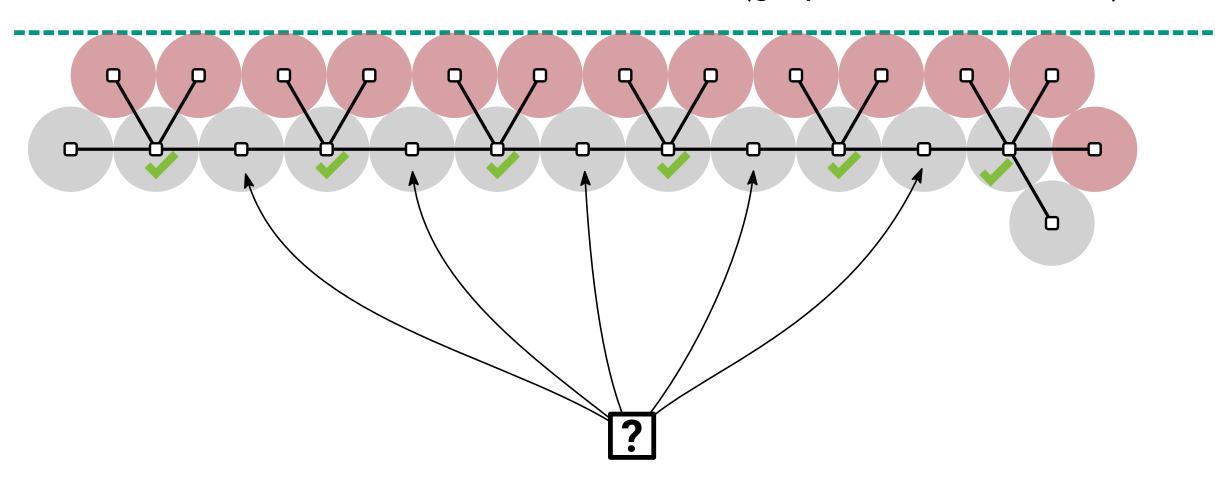


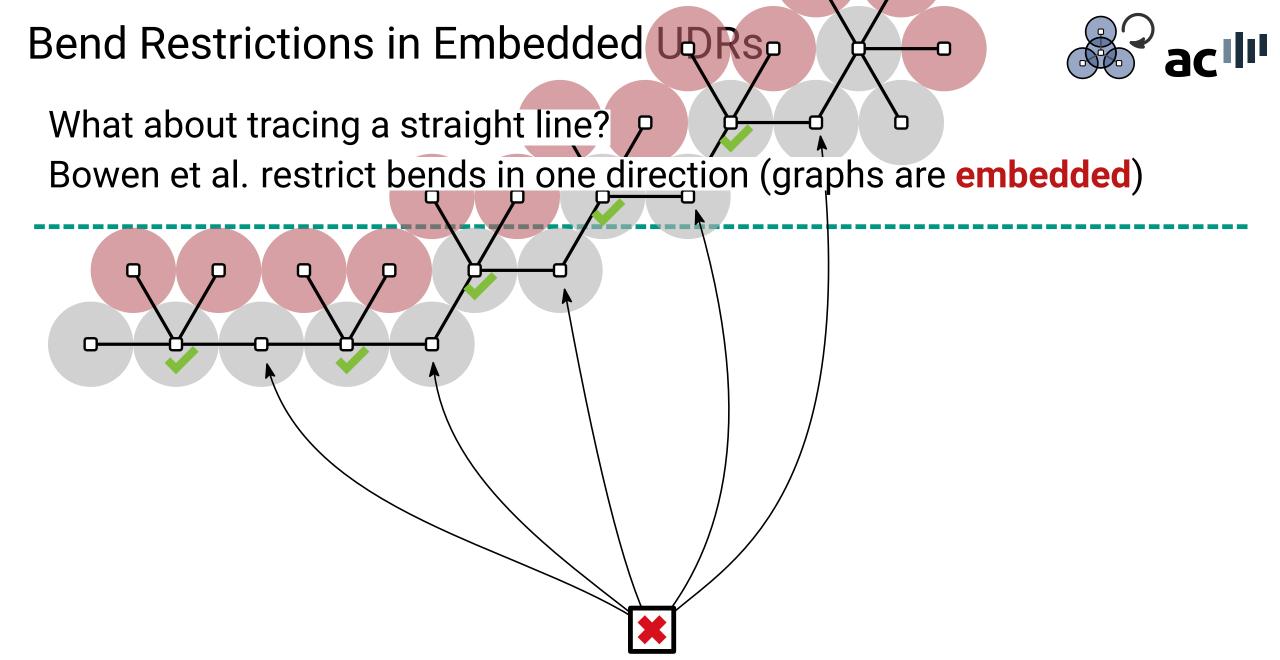
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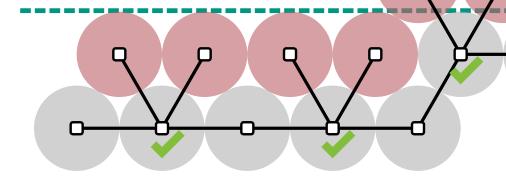


We need additional restricting neighbors

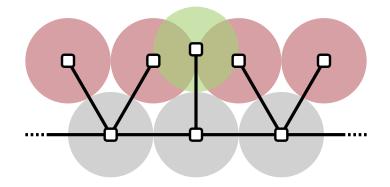


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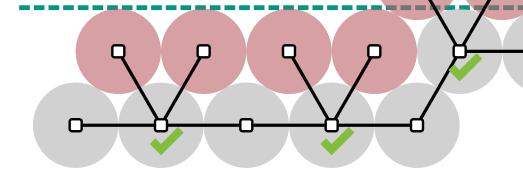
We can force nearly complete overlap with a single disk at these nodes



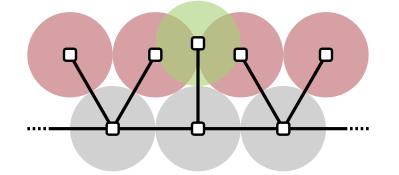


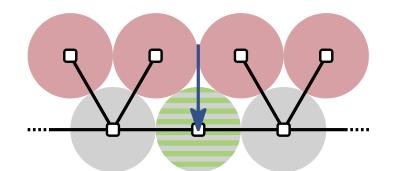
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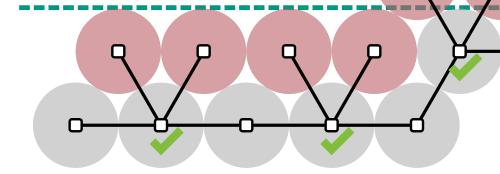




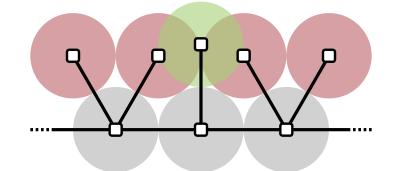


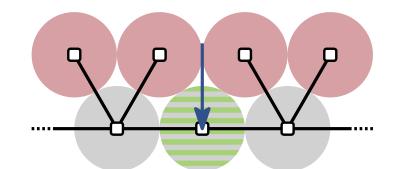
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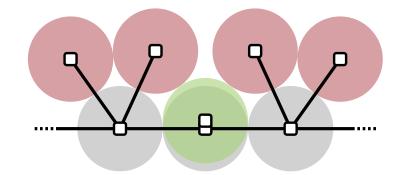
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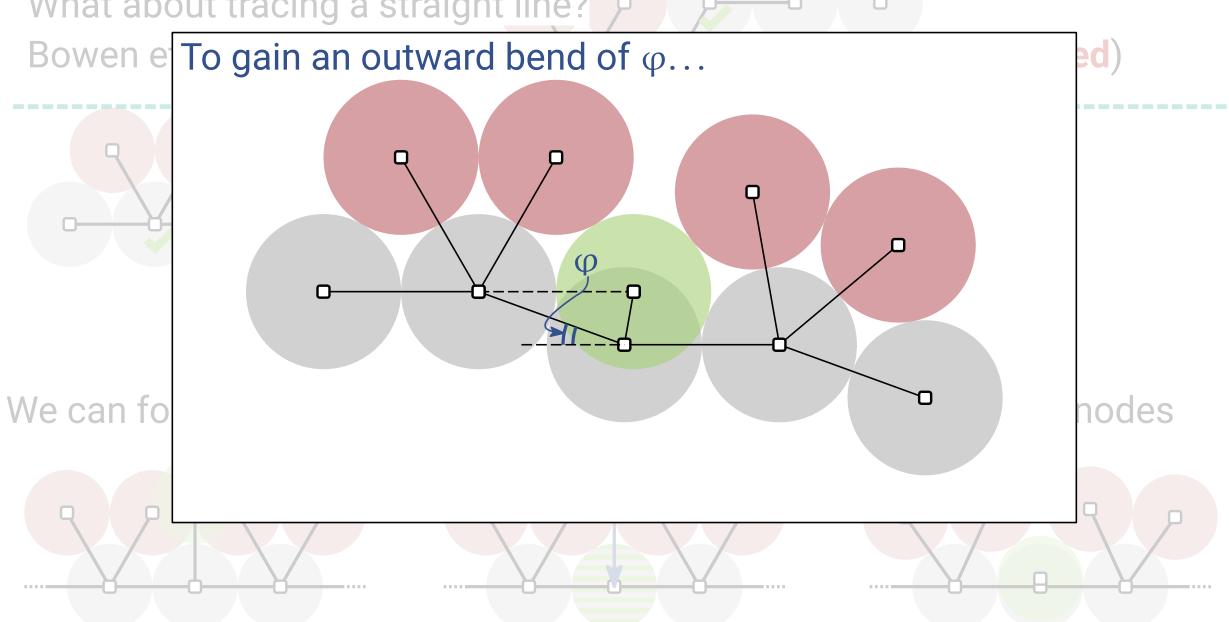
Bend Restrictions in Embedded UDRsp What about tracing a straight line? P Bowen e

nodes

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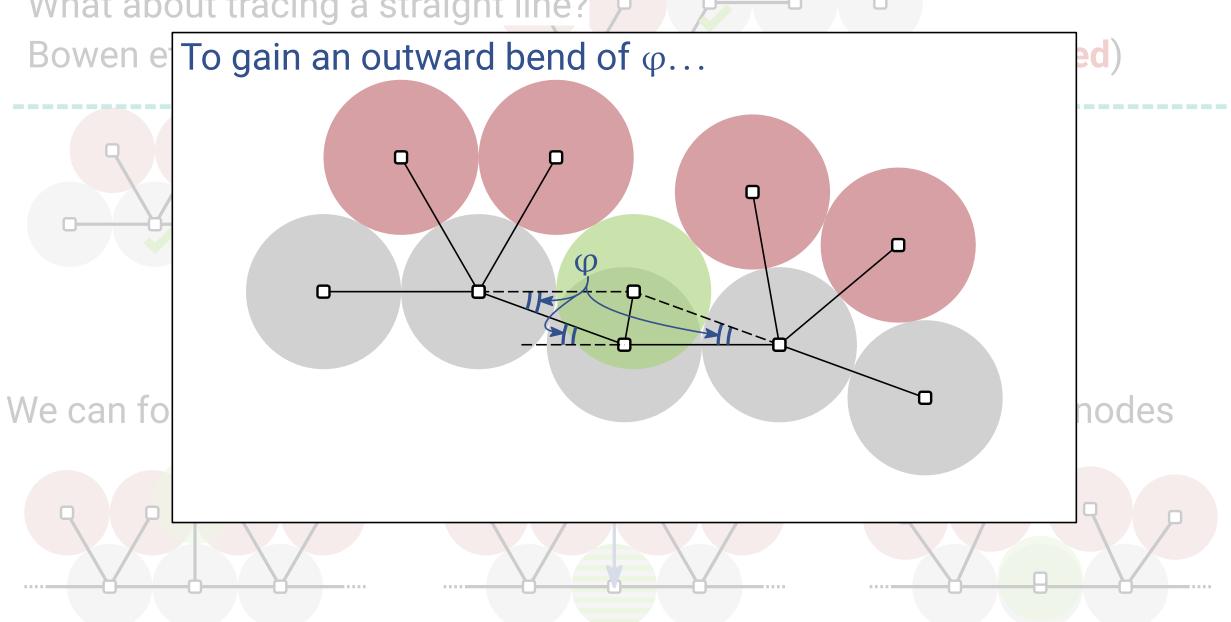
Bend Restrictions in Embedded UDRs What about tracing a straight line? Bowen e To gain an outward bend of ϕ ...





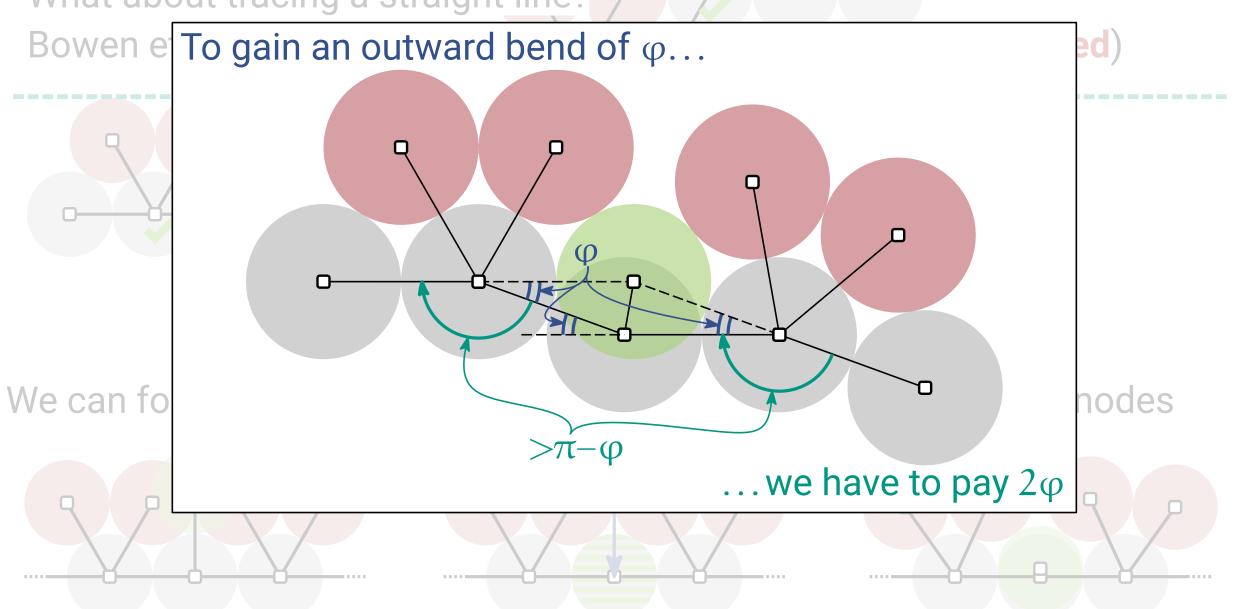
Bend Restrictions in Embedded UDRsp What about tracing a straight line?





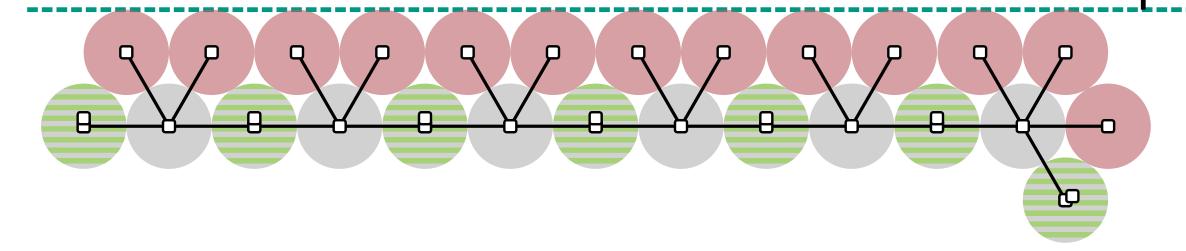


What about tracing a straight line? P

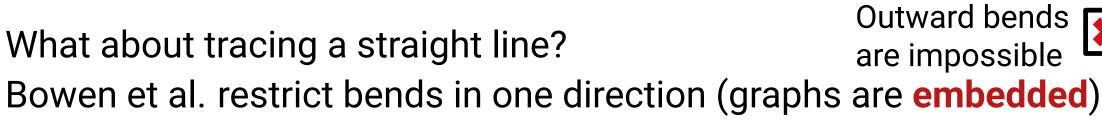


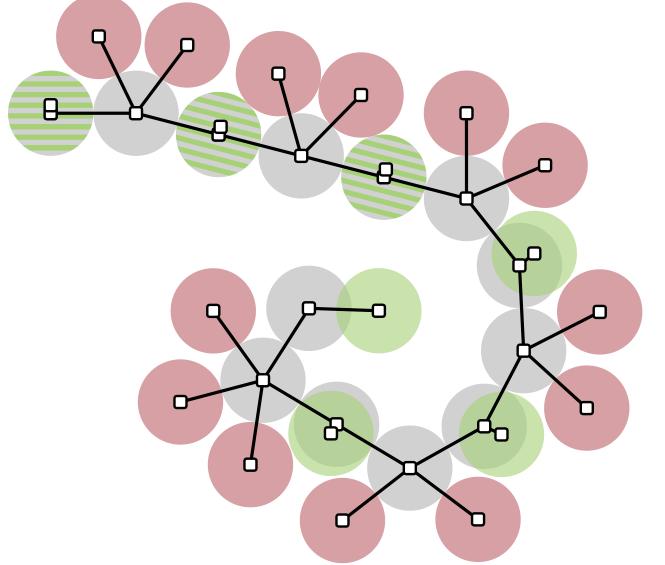
Outward bends are impossible

What about tracing a straight line?



What about tracing a straight line?

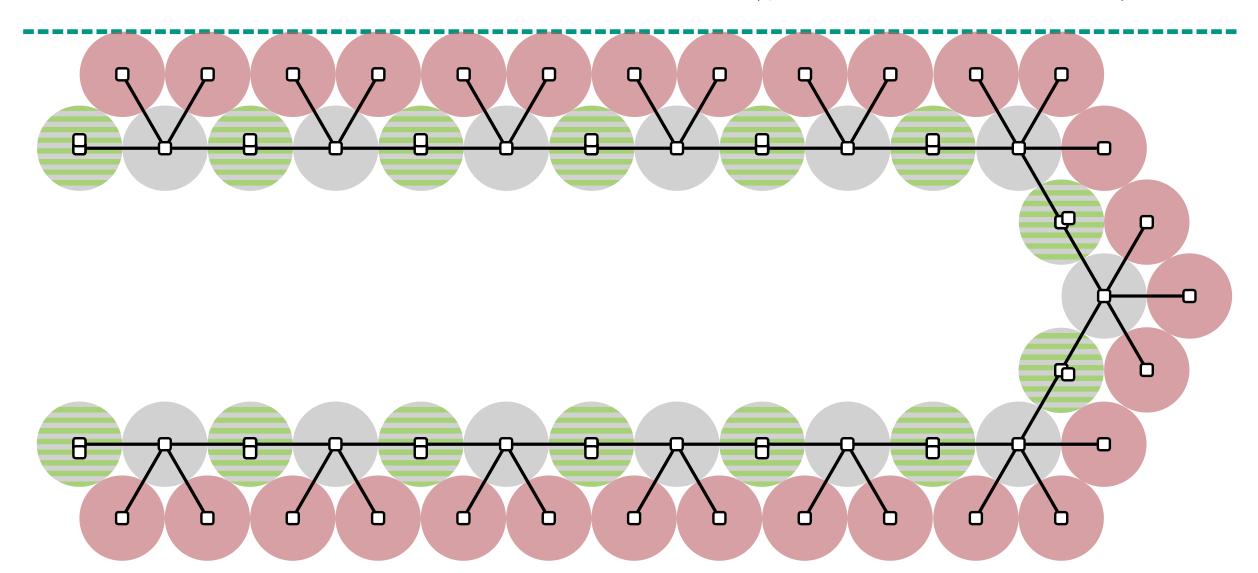




How do we treat inward flexibility?

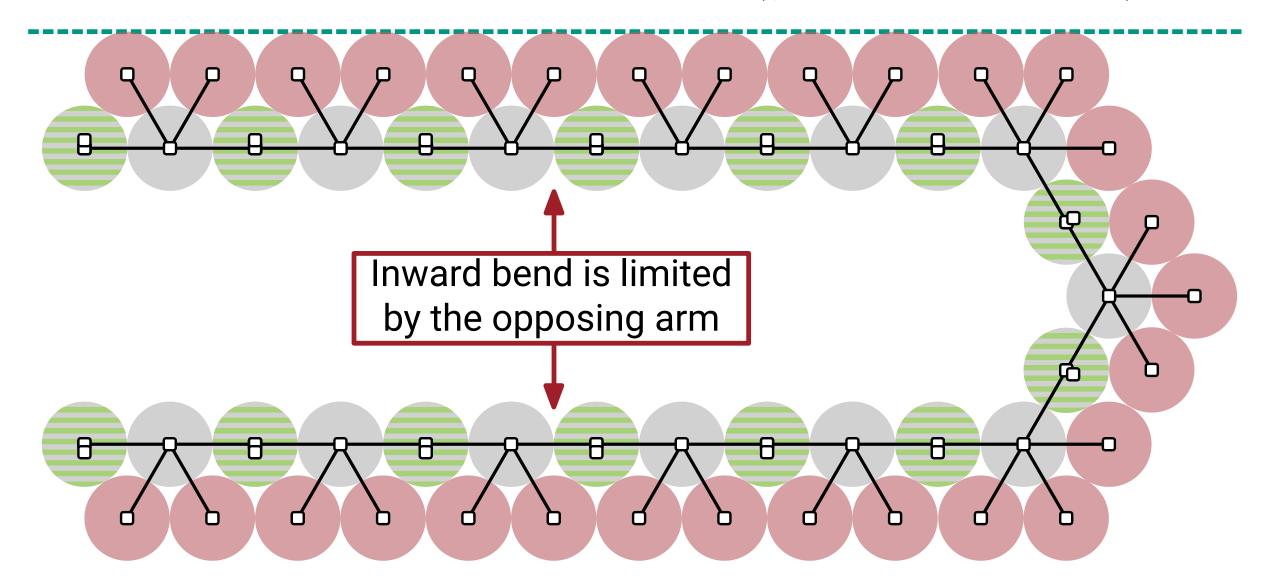


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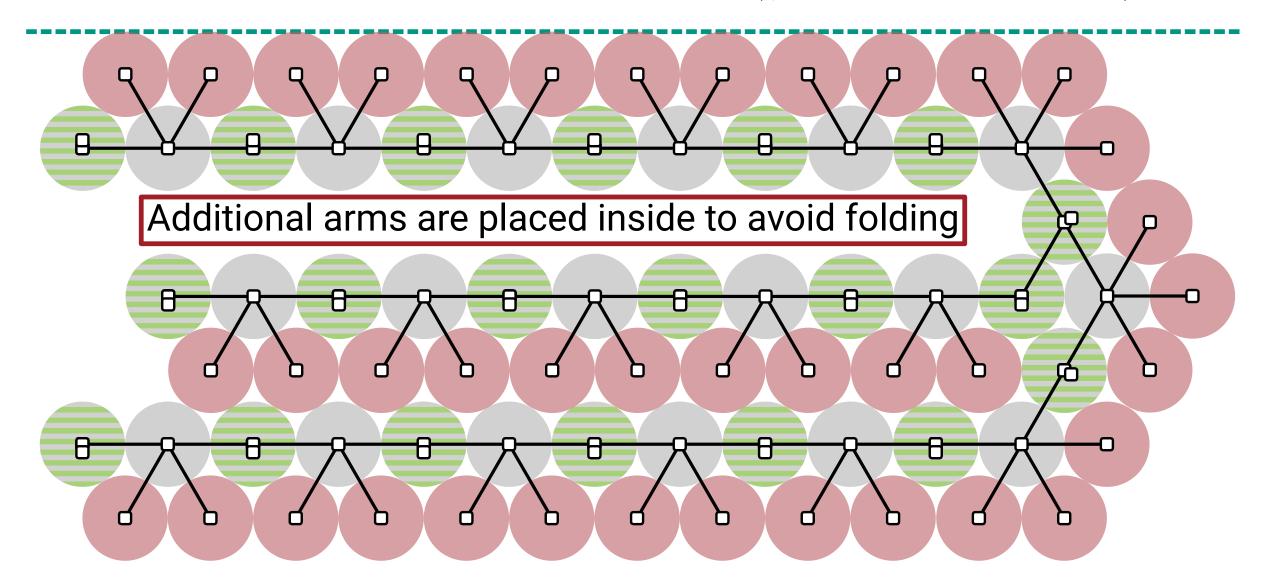


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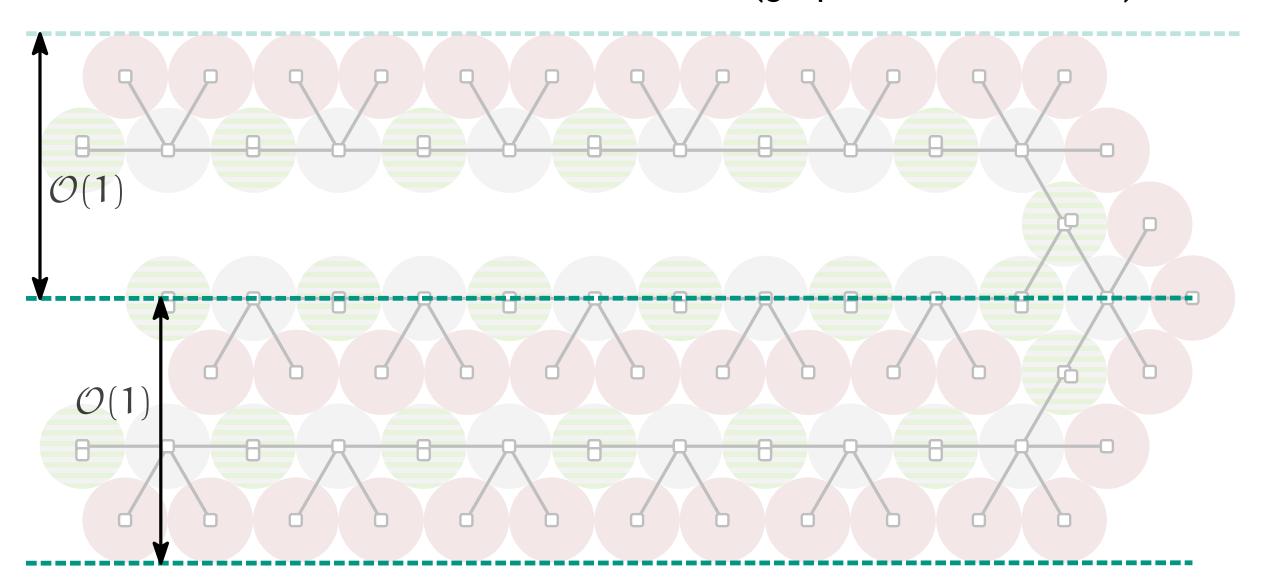


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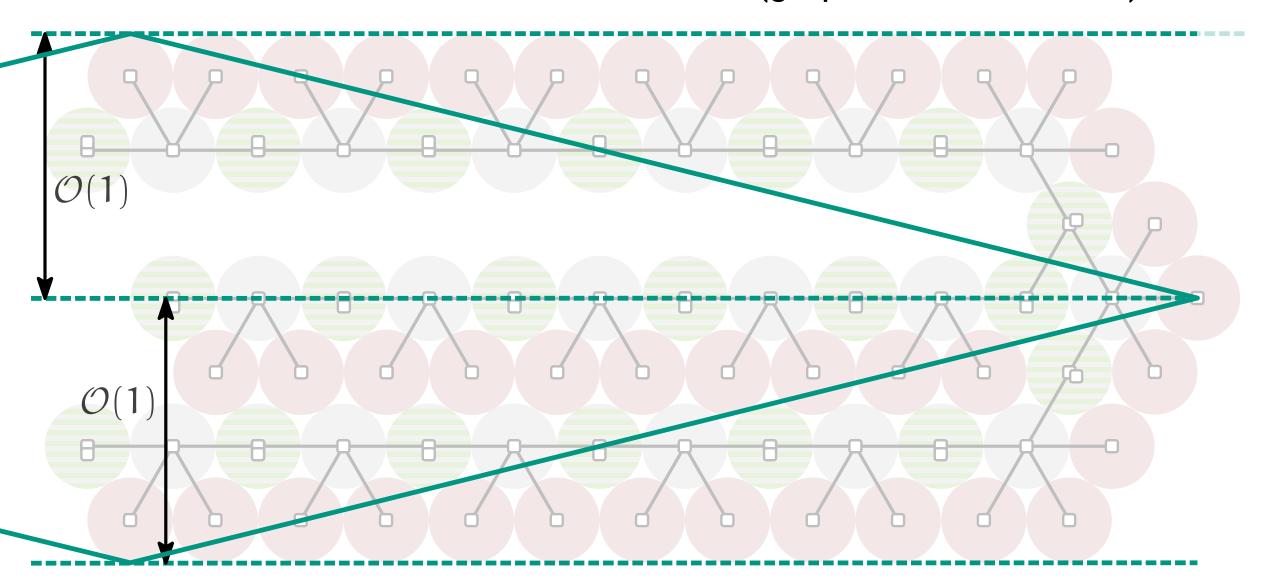


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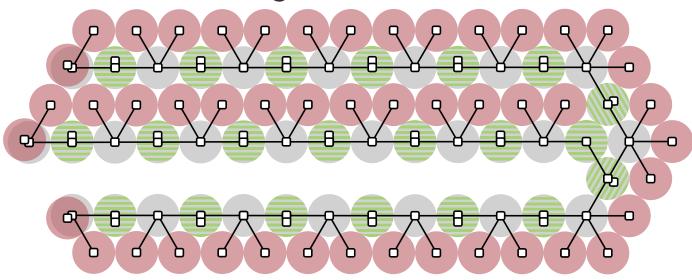


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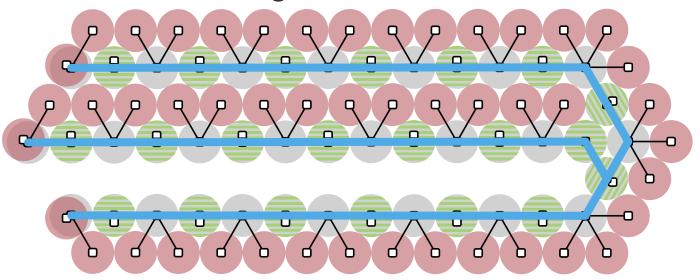


 $\mathcal{O}(1)$ -stable approximation of a long thin rhombus

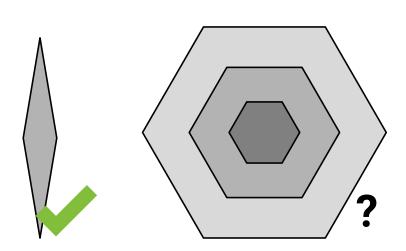




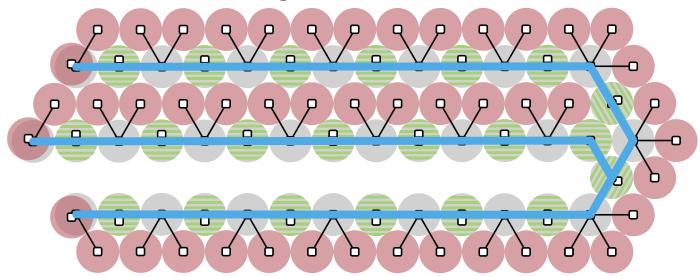
 $\mathcal{O}(1)$ -stable approximation of a long thin rhombus



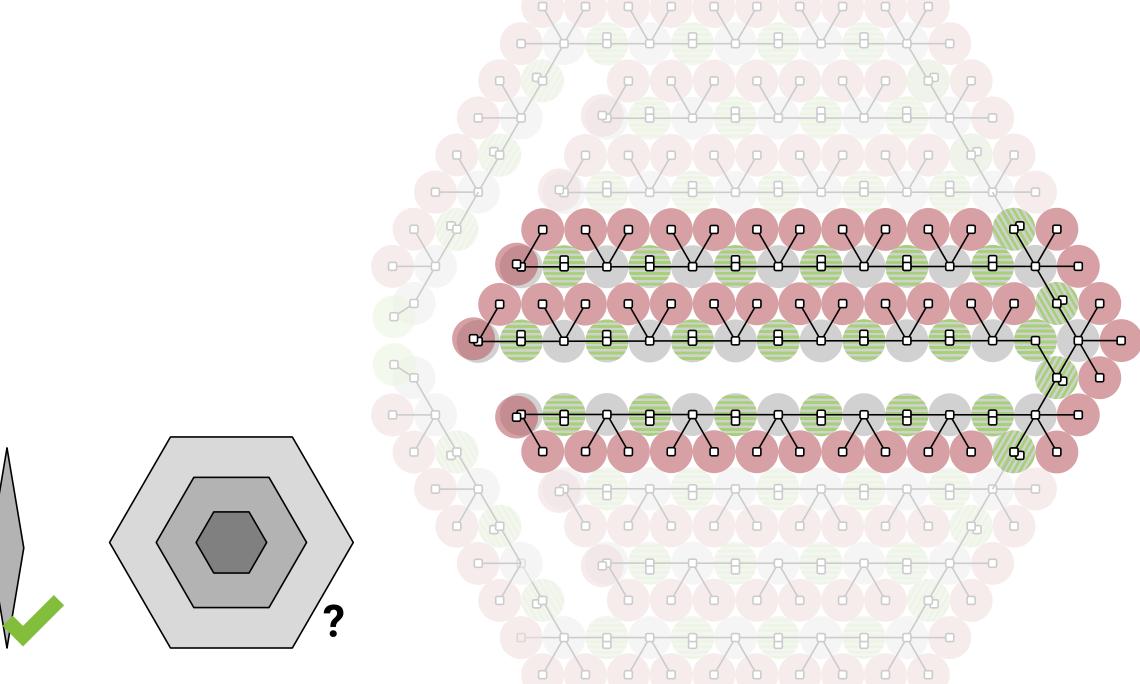


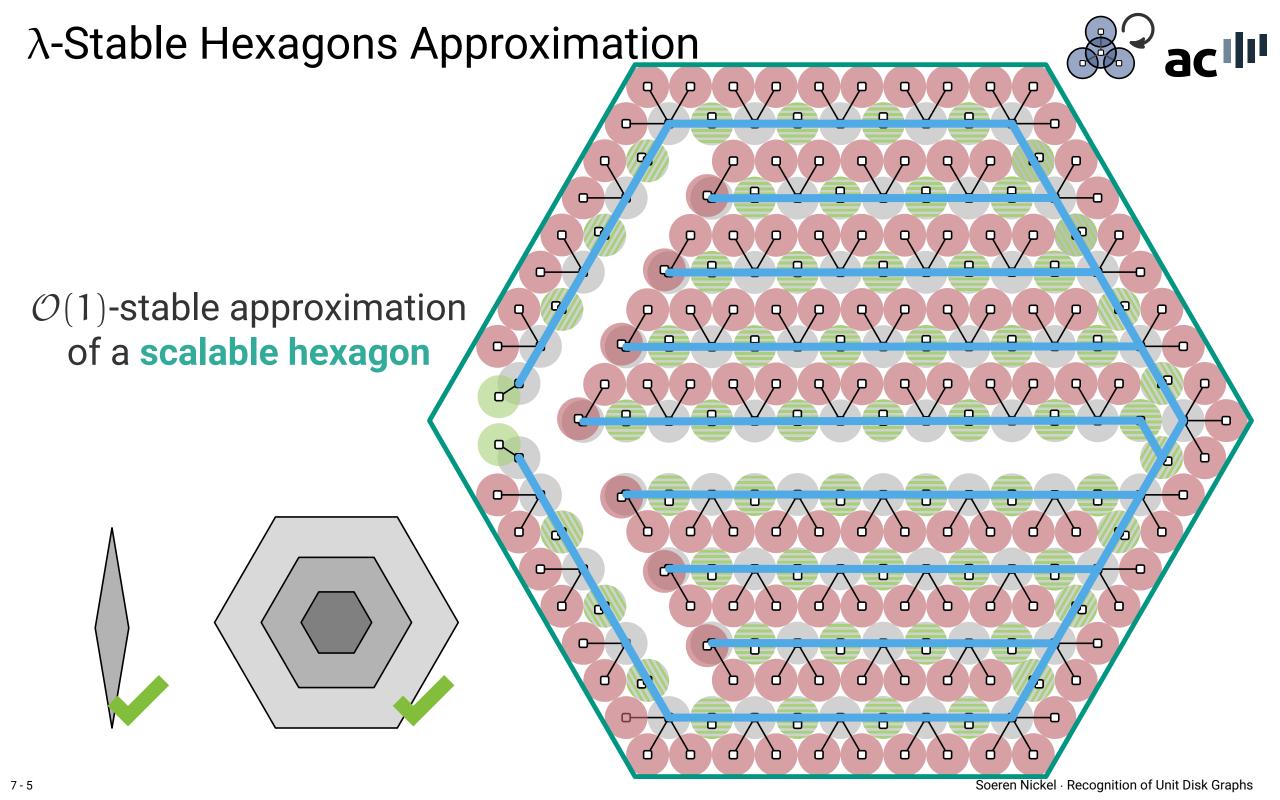


 $\mathcal{O}(1)$ -stable approximation of a long thin rhombus









Hardness Results and Outerplanarity



Theorem 1. Recognizing unit disk graphs is NP-hard for embedded trees.

Hardness Results and Outerplanarity

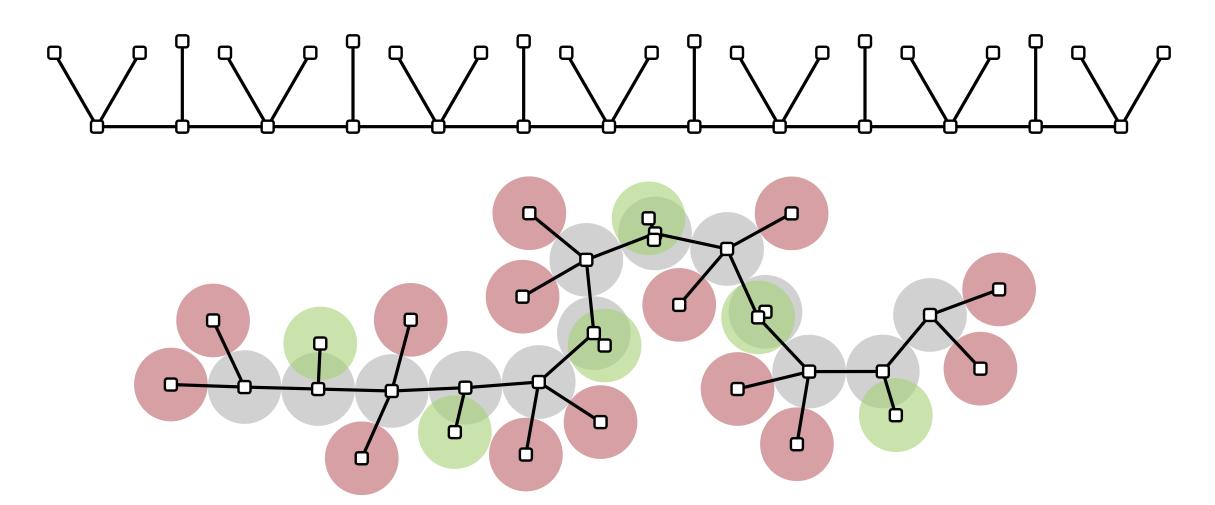


Theorem 1. Recognizing unit disk graphs is NP-hard for embedded trees.

What about **non-embedded** graphs?



Theorem 1. Recognizing unit disk graphs is NP-hard for embedded trees. What about **non-embedded** graphs?

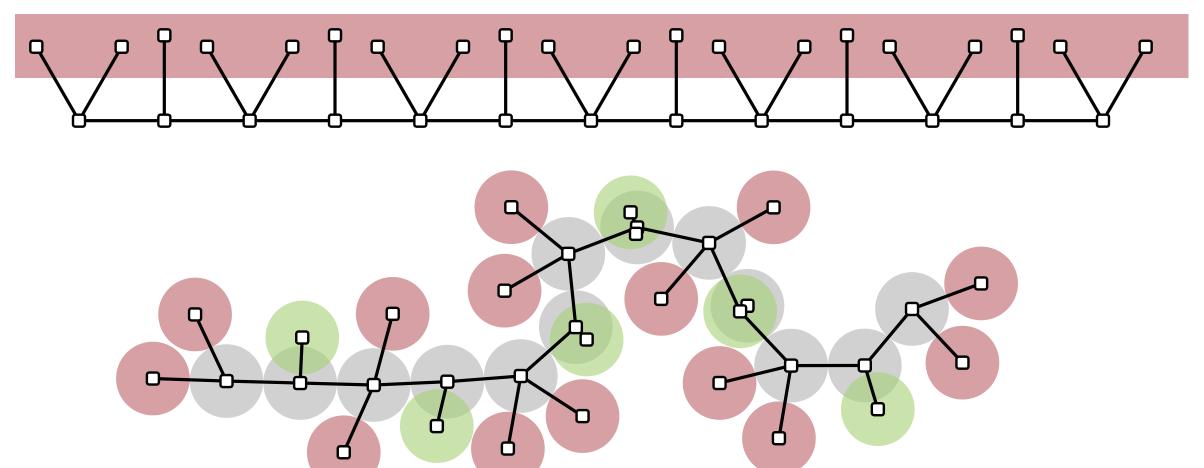




Theorem 1. Recognizing unit disk graphs is NP-hard for embedded trees.

What about non-embedded graphs?

These need to be forced on one side

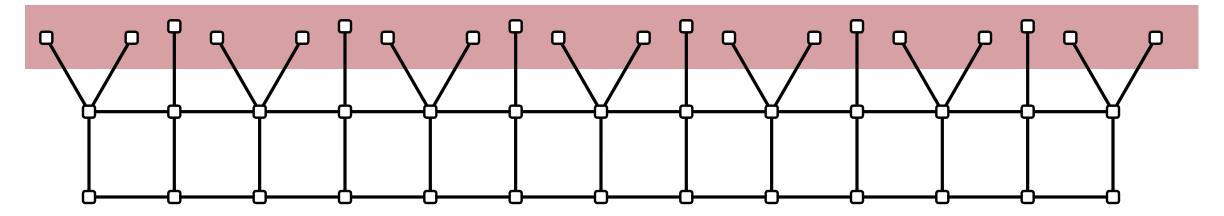




Theorem 1. Recognizing unit disk graphs is NP-hard for embedded trees.

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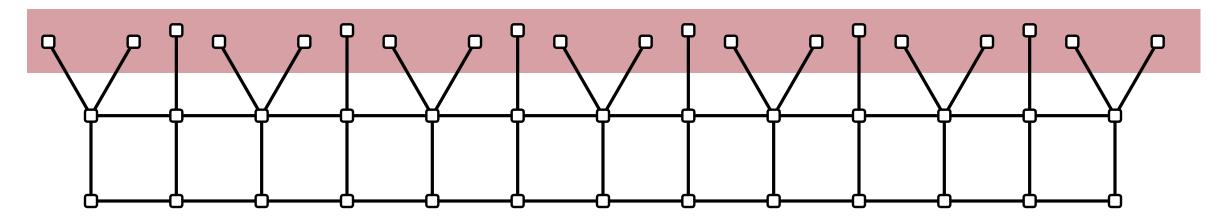




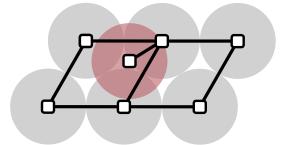
Theorem 1. Recognizing unit disk graphs is NP-hard for embedded trees.

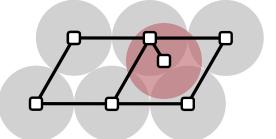
What about non-embedded graphs?

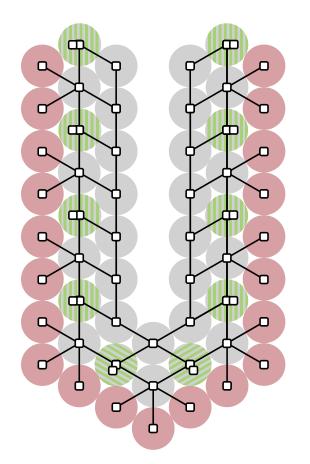
These need to be forced on one side

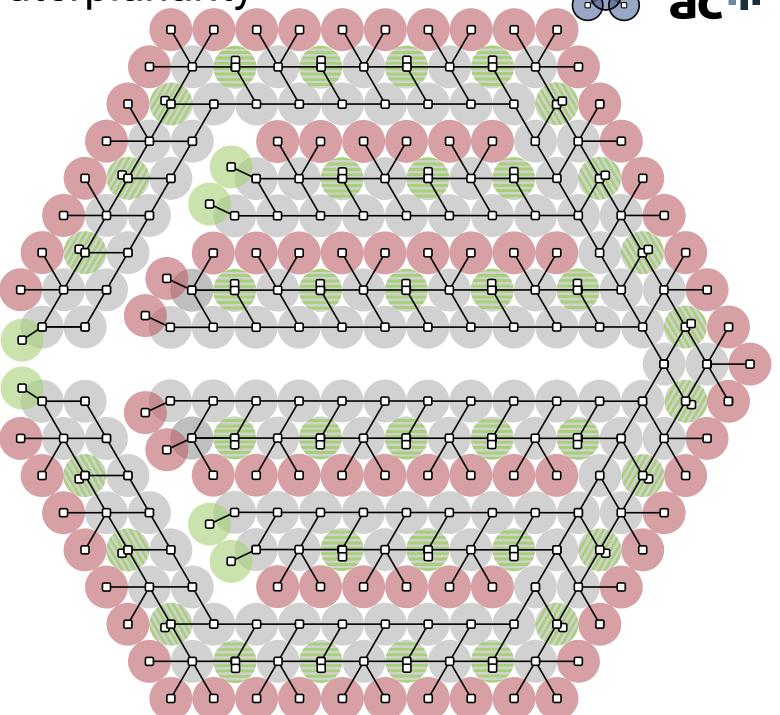


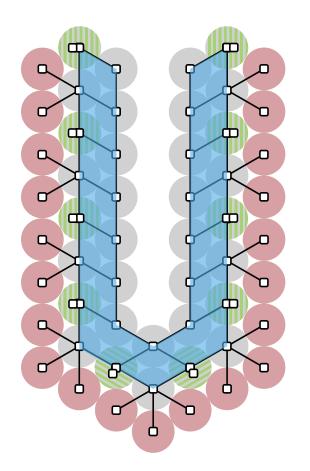
Outer neighbours can not be placed on the inside.

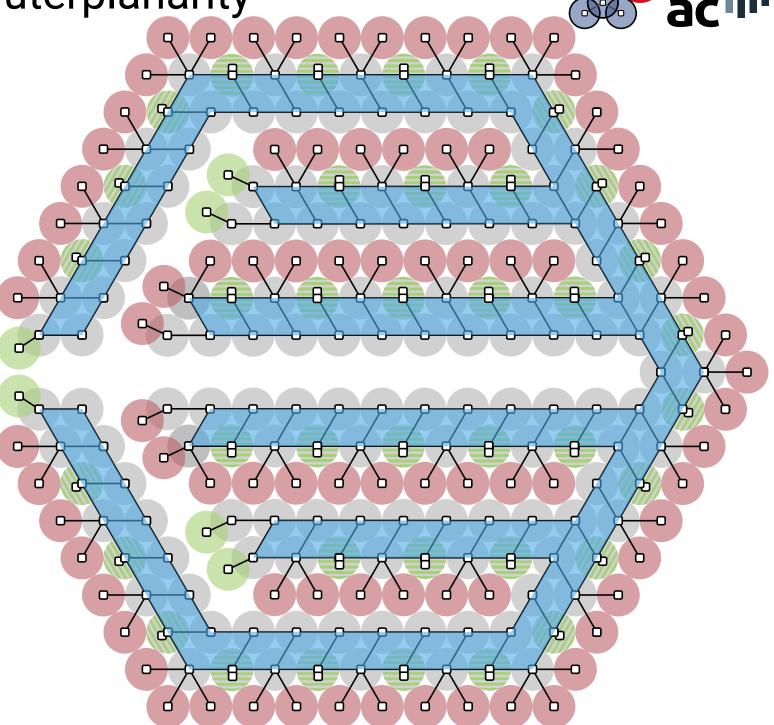










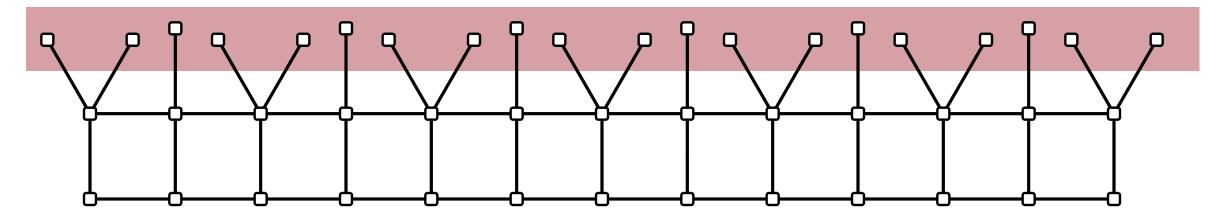




Theorem 1. Recognizing unit disk graphs is NP-hard for embedded trees.

What about **non-embedded** graphs?

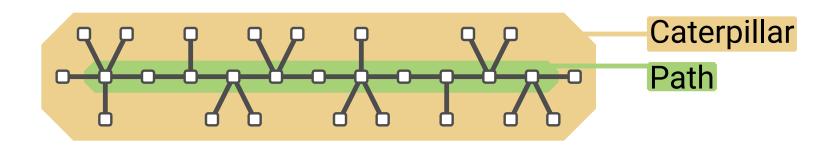
These need to be forced on one side



Theorem 2. Recognizing unit disk graphs is NP-hard for non-embedded outerplanar graphs.



Caterpillars without embedding





Caterpillars without embedding

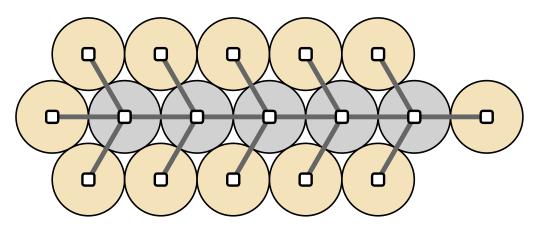
A caterpillar C





Caterpillars without embedding

A caterpillar C



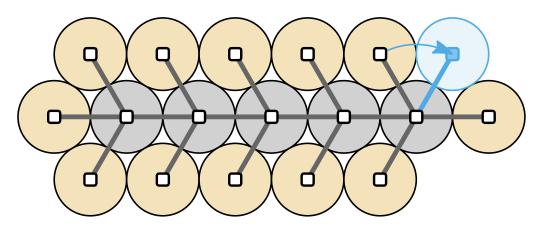
4-regular spines





Caterpillars without embedding

A caterpillar C

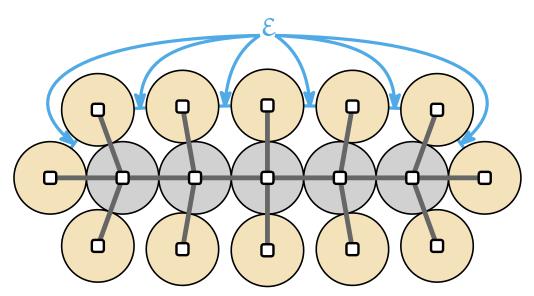


4-regular spines





Caterpillars without embedding



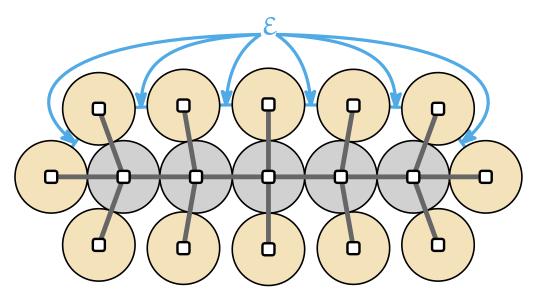
4-regular spines are trivial!

A caterpillar C



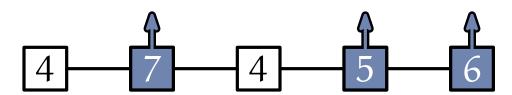


Caterpillars without embedding



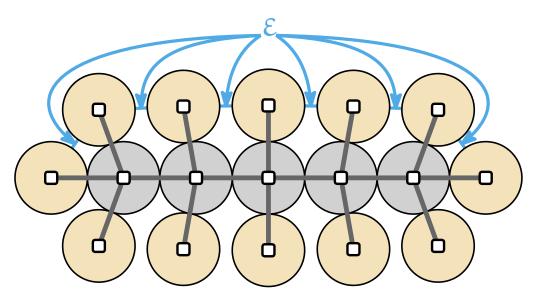
4-regular spines are trivial!



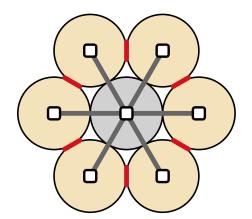




Caterpillars without embedding

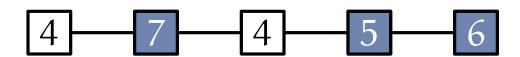


4-regular spines are trivial!



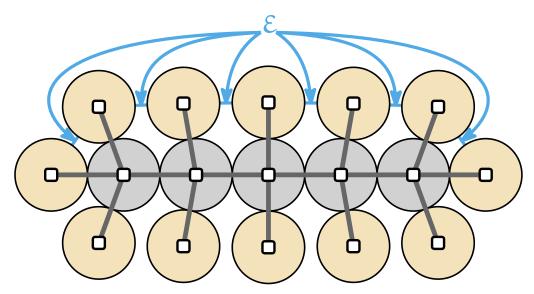
Degree 6 vertices are **impossible**!

A caterpillar C

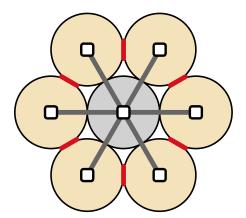




Caterpillars without embedding

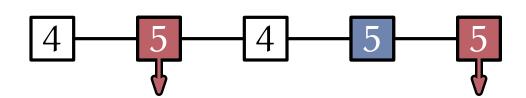


4-regular spines are trivial!



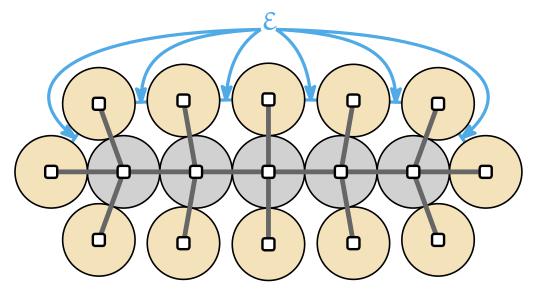
Degree 6 vertices are impossible!

A caterpillar C

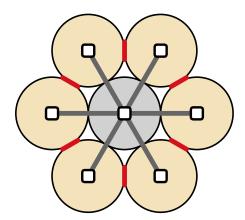




Caterpillars without embedding



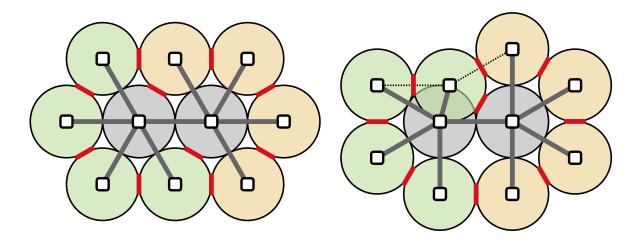
4-regular spines are trivial!



Degree 6 vertices are **impossible**!

A caterpillar C

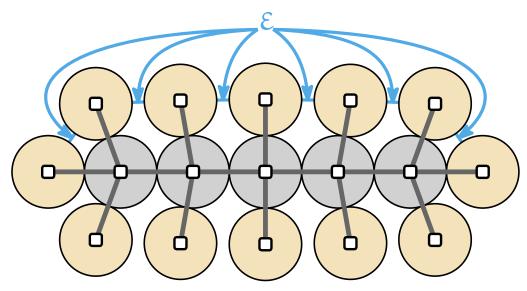




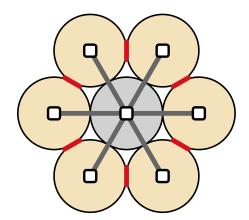
Consecutive degree 5 vertices are impossible!



Caterpillars without embedding

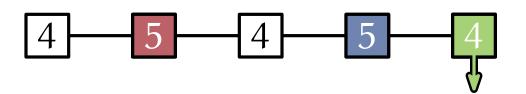


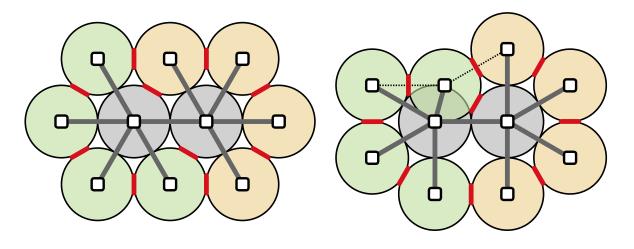
4-regular spines are trivial!



Degree 6 vertices are **impossible**!

A caterpillar C

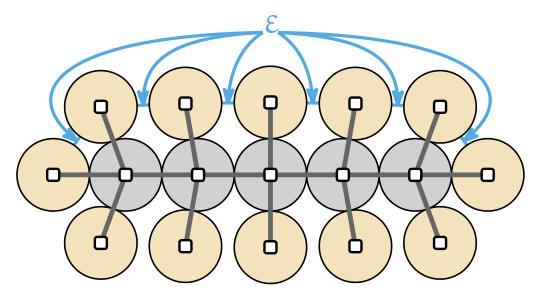




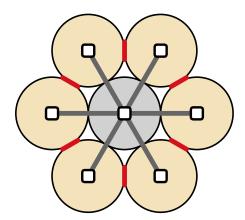
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Caterpillars without embedding

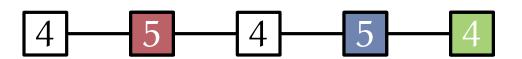


4-regular spines are trivial!

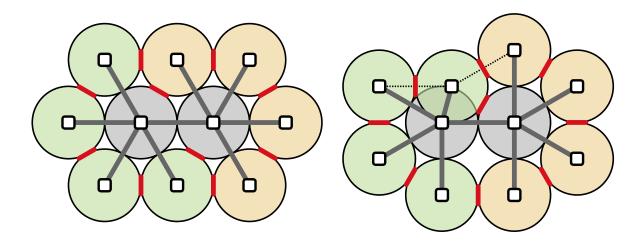


Degree 6 vertices are **impossible**!

A caterpillar C



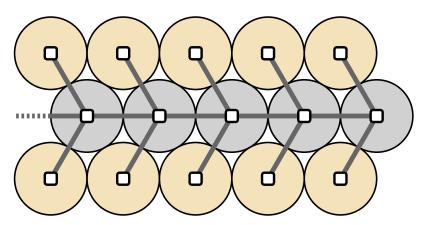
Can we realize $C = \dots 5454545\dots$?



Consecutive degree 5 vertices are impossible!

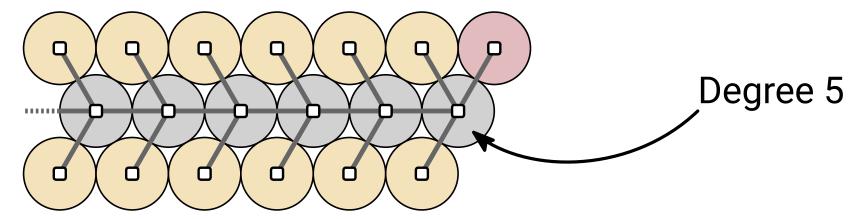


Caterpillars without embedding





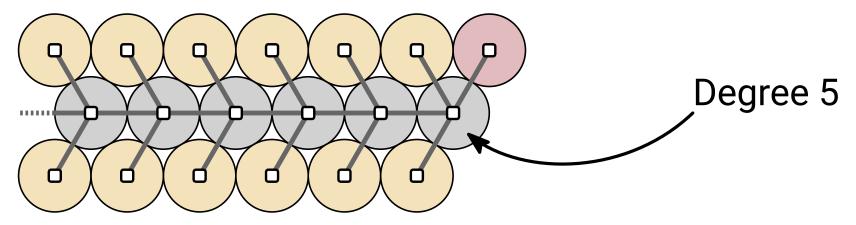
Caterpillars without embedding





Caterpillars without embedding

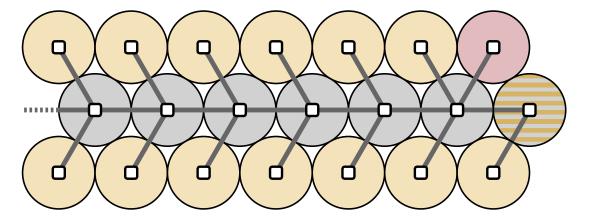
Backward position, as long as degree is ≤ 4



If the next vertex has degree 5, we abort.



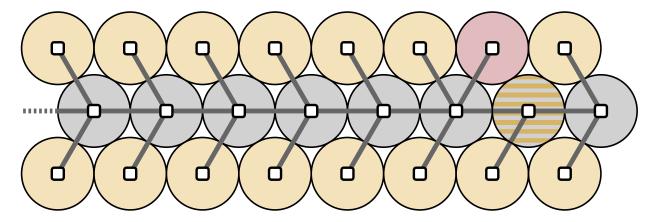
Caterpillars without embedding



- If the next vertex has degree 5, we abort.
- Otherwise we can return to backward position.



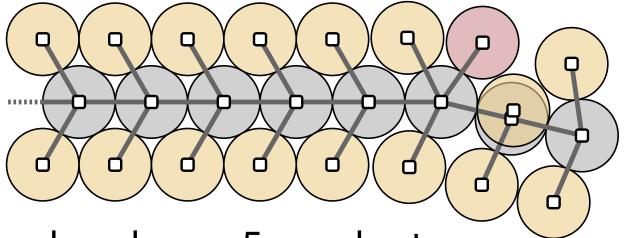
Caterpillars without embedding



- If the next vertex has degree 5, we abort.
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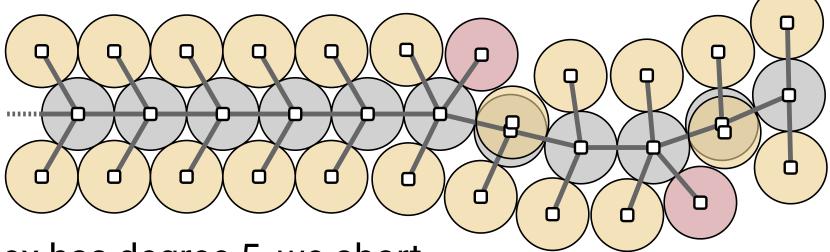
Caterpillars without embedding



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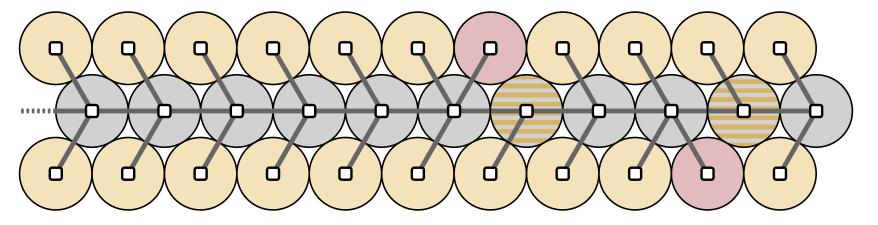
Caterpillars without embedding



- If the next vertex has degree 5, we abort.
- Otherwise we can return to backward position.
- At the end, use available slack again to avoid wrong adjacencies



Caterpillars without embedding

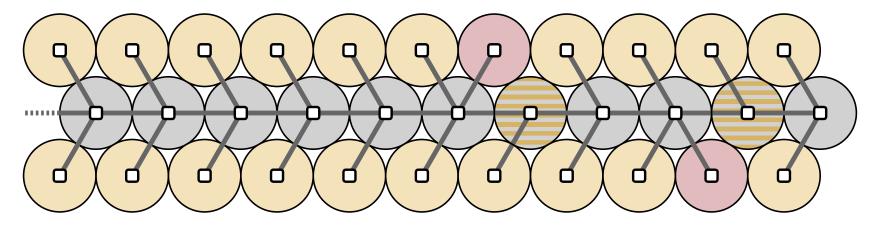


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Caterpillars without embedding

Backward position, as long as degree is ≤ 4



- If the next vertex has degree 5, we abort.
- Otherwise we can return to backward position.
- At the end, use available slack again to avoid wrong adjacencies

Theorem 3. Caterpillar graphs admit a UDR, if and only if they neither contain vertices of degree 6 or higher nor two adjacent vertices of degree 5.



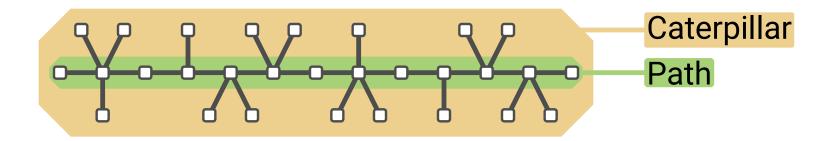


Recognizing unit disk graphs is NP-hard for embedded trees and outerplanar graphs (non-embedded).

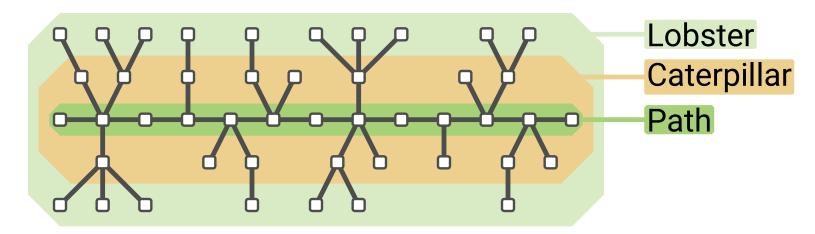


- Recognizing unit disk graphs is NP-hard for embedded trees and outerplanar graphs (non-embedded).
- Recognition can be done in linear time for caterpillars (non-embedded)

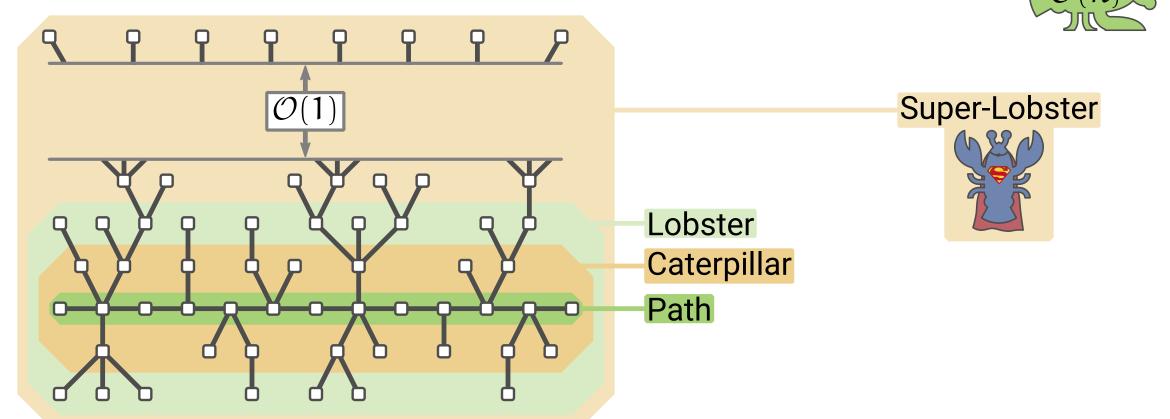
- Recognizing unit disk graphs is NP-hard for embedded trees and outerplanar graphs (non-embedded).
- Recognition can be done in linear time for caterpillars (non-embedded)
- Computer-assisted proof: lobster graphs on a grid can be represented as a x-monotone weak UDC and recognized in linear time



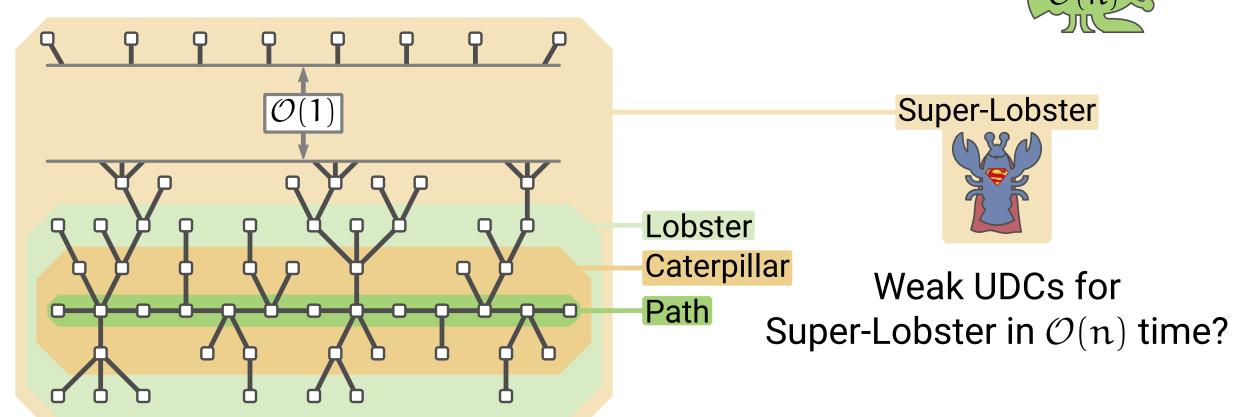
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- Recognizing unit disk graphs is NP-hard for embedded trees and outerplanar graphs (non-embedded).
- Recognition can be done in linear time for caterpillars (non-embedded)
- **Computer-assisted** proof: **lobster graphs** on a **grid** can be represented as a x-monotone *weak* UDC and recognized in **linear time**



- Recognizing unit disk graphs is NP-hard for embedded trees and outerplanar graphs (non-embedded).
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Open questions:		tree	caterpillar
С	ontact graphs	?	$\mathcal{O}(\mathfrak{n})$
	→ embedded	NP-hard	?
disk graphs		?	$\mathcal{O}(\mathfrak{n})$
	→ embedded	NP-hard	?



Weak UDCs for Super-Lobster in O(n) time?